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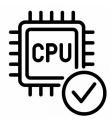
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Symbolic Methods

- Background
 - Applications
 - Verify complex systems.
 - Theorem proving.
 - Constraint solving in optimization.
 - Symbolic expressions effectively encode large sets of states and transitions.









Symbolic Methods

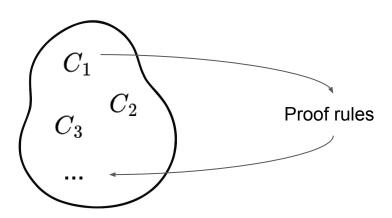
- Fundamental techniques in automated reasoning.
 - Deductive reasoning, term rewriting, constraint solving, etc.
- Approaches
 - Automated Theorem Provers (ATPs).
 - Boolean Satisfiability (SAT)/Satisfiability Modulo Theories (SMT) solvers.
 - Constrained Horn Clause (CHC) solvers.

Motivating Examples (Program Verification)

$$egin{aligned} x = _ & \longrightarrow C_1: \ L_1(x) & \leftarrow true \ ext{while}(x > 0) \{ & \longrightarrow C_2: \ L_2(x) & \leftarrow L_1(x) \land x > 0 \ x = x - 1 & \longrightarrow C_3: \ L_1(x') & \leftarrow L_2(x) \land x' = x - 1 \ \} & C_4: \ L_3(x) & \leftarrow L_1(x) \land x \leq 0 \ ext{assert}(x
eq 0) & \longrightarrow C_5: \ false & \leftarrow L_3(x) \land x = 0 \end{aligned}$$

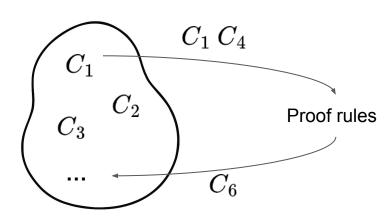
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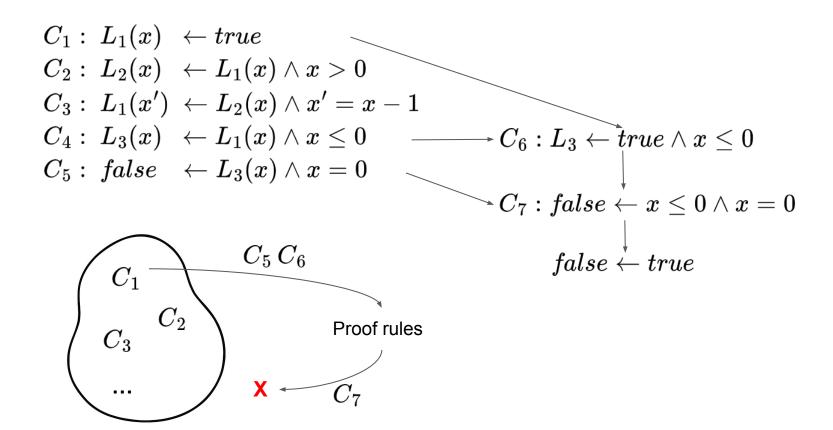


Motivating Examples

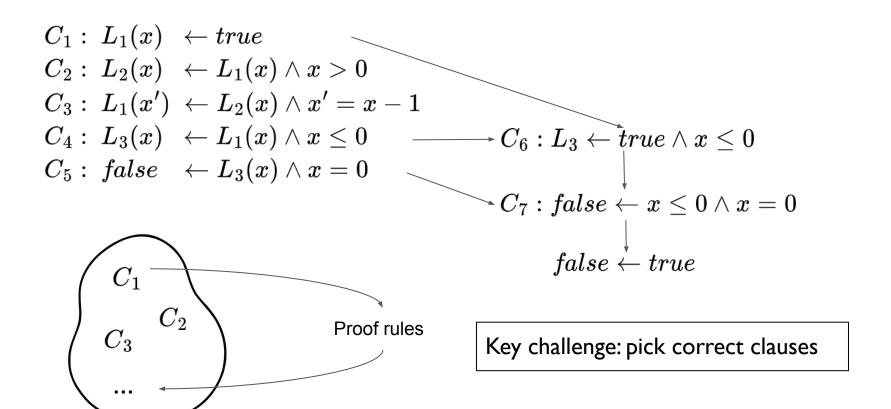
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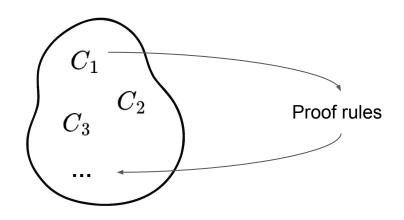
Motivating Examples



Motivating Examples

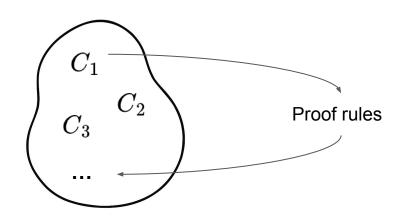


- Need heuristics to pick clauses
 - Predefined features [1]:
 - Clause age.
 - Clause size.
 - Relevance to the proof goal.

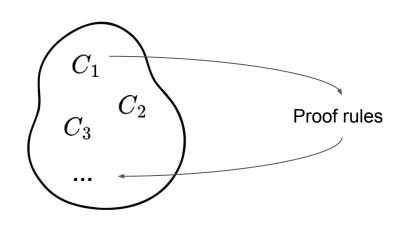


[1] S. Schulz and M. Möhrmann, "Performance of clause selection heuristics for saturation-based theorem proving," 2016.

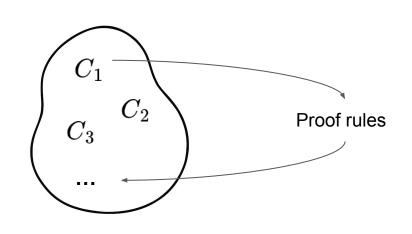
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 - Data-driven.
- Deepire [2] uses recursive neural networks to classify the clauses based on their derivation history.

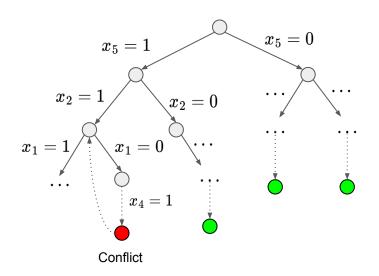


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- ENIGMA [3] uses both Gradient Boosting Decision Trees (GBDTs) and Graph Neural Networks (GNNs) to select the clauses.



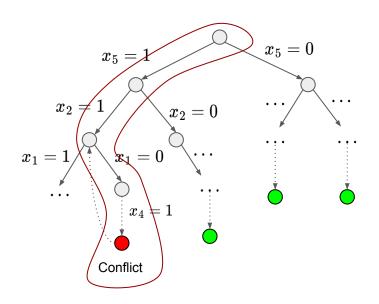
Clauses

$$(x_1ee x_4)\wedge \ (x_3ee ar x_4ee ar x_5)\wedge \ (ar x_3ee ar x_2ee ar x_4)$$



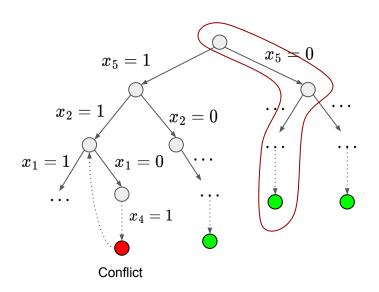
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Clauses

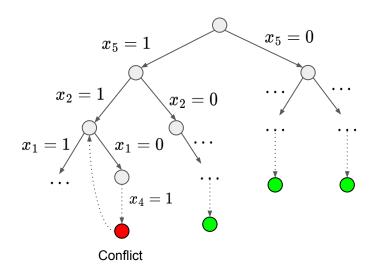
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Clauses

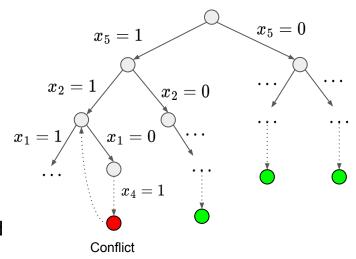
$$(x_1ee x_4)\wedge \ (x_3ee ar x_4ee ar x_5)\wedge \ (ar x_3ee ar x_2ee ar x_4)$$

- SAT Solver
 - Selecting an variable for branching.
 - Which branch to go.
 - Deciding when to restart.

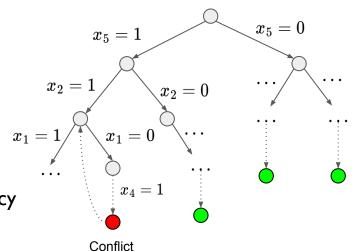


Key challenge: find the correct path

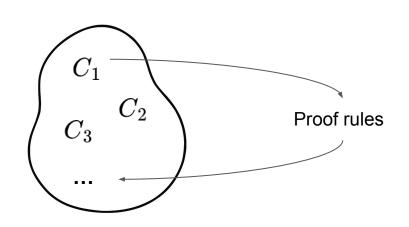
- Conflict-Driven Clause Learning (CDCL)-based SAT solving
 - Selecting an variable for branching.
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- □ NeuroSAT [4] periodically resets Exponential Variable State-Independent Decaying Sum (EVSIDS) scores based on the predictions of a message-passing neural network.

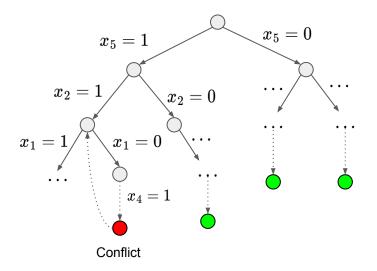


- Conflict-Driven Clause Learning (CDCL)-based SAT solving
 - Selecting an variable for branching.
 - Which branch to go.
 - Deciding when to restart.
- Authors in [5] propose a data-driving-based restart policy by analyzing the history of previously learned clauses.

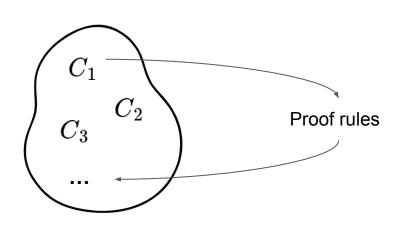


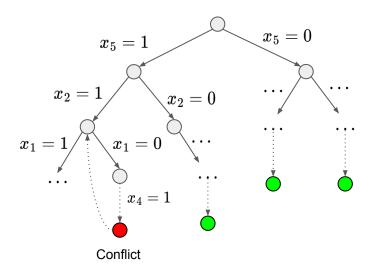
[5] J. H. Liang, C. Oh, M. Mathew, C. Thomas, C. Li, and V. Ganesh, "Machine learning-based restart policy for CDCL SAT solvers," 2018.



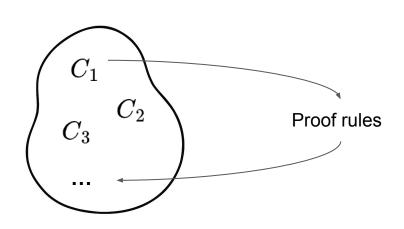


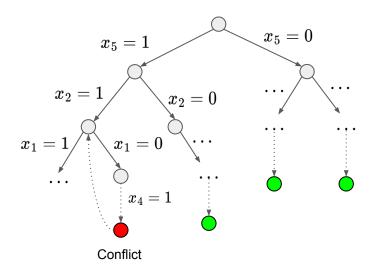
 Nearly all approaches of symbolic methods incorporate multiple heuristic-driven decision processes.



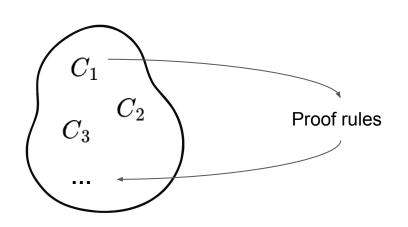


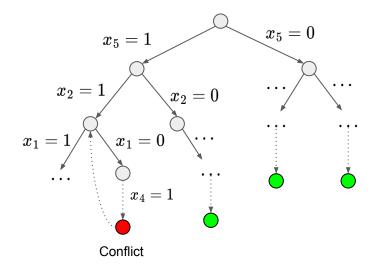
- Nearly all approaches of symbolic methods incorporate multiple heuristic-driven decision processes.
- Deep-learning guided heuristics show promising results.





- □ Nearly all approaches of symbolic methods incorporate multiple heuristic-driven decision processes.
- Deep-learning guided heuristics show promising results.
- □ Each deep-learning guided heuristic is highly customized.



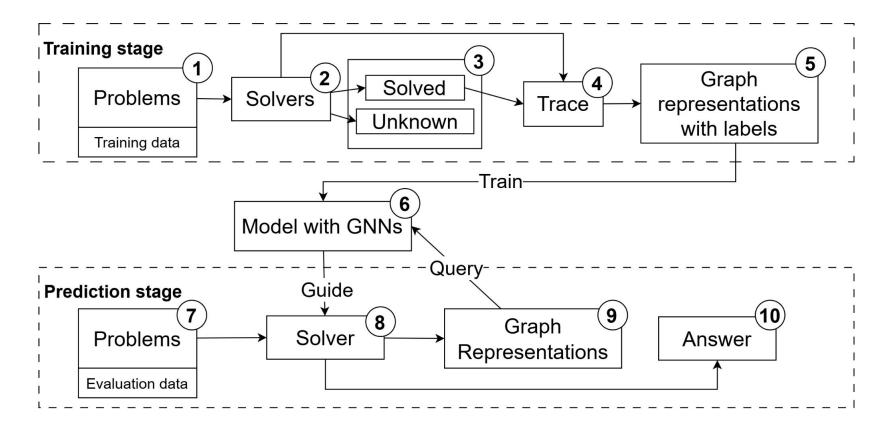


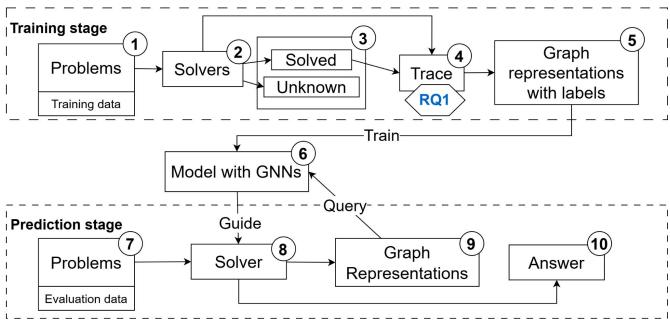
Can we generalize the deep-learning based guiding systematically, so that the decision processes can be improved within a unified framework?

List of papers

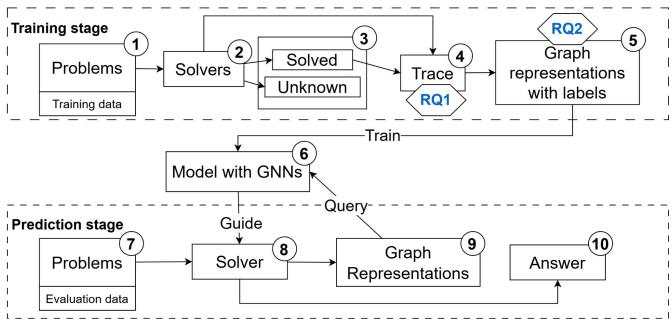
- **I Exploring Representation of Horn Clauses Using GNNs.** Chencheng Liang, Philipp Rümmer, Marc Brockschmidt. *In Proceedings of 8th Workshop on Practical Aspects of Automated Reasoning (PAAR)*, 2022.
- **II Boosting Constrained Horn Solving by Unsat Core Learning.** Parosh Aziz Abdulla, Chencheng Liang, Philipp Rümmer. *In Proceedings of 25th International Conference on Verification, Model Checking, and Abstract Interpretation (VMCAI)*, 2024.
- III Guiding Word Equation Solving Using Graph Neural Networks. Parosh Aziz Abdulla, Mohamed Faouzi Atig, Julie Cailler, Chencheng Liang, Philipp Rümmer. In Proceedings of 22nd International Symposium on Automated Technology for Verification and Analysis (ATVA), 2024.
- **IV** When GNNs Met a Word Equations Solver: Learning to Rank Equations. Parosh Aziz Abdulla, Mohamed Faouzi Atig, Julie Cailler, Chencheng Liang, Philipp Rümmer. *Under Submission*, 2025.
- Deep-learning based framework
- Instance I: Constrained Horn clauses (CHCs) solving (paper I and II)
- Instance 2: Word equation solving (paper III and IV)
- Conclusion and future works

Deep Learning-Based Framework (Syntactic Structure)

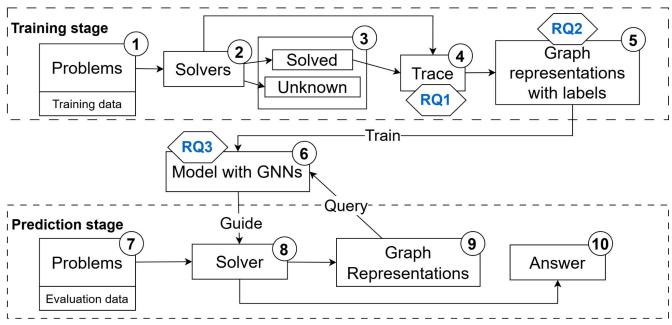




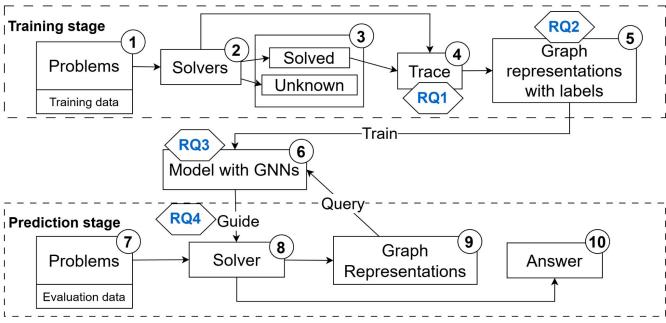
RQI:What are good encodings of symbolic decision processes as training tasks?



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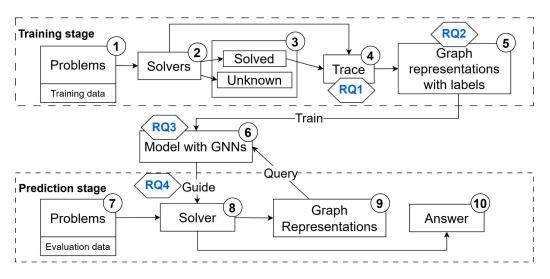


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- RQ2:What is the most effective format for representing formulas in deep learning?
- □ RQ3:Which deep learning technique is best suited for feature extraction from formulas?
- RQ4:What are the methods for integrating the trained model into algorithms?

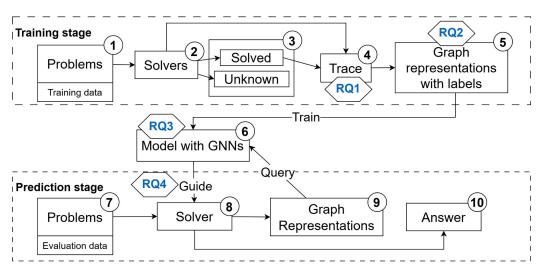
Deep Learning-Based Framework (Instances)



Guide Constrained Horn Clauses (CHCs) solving (Paper I and II)

$$egin{aligned} F(0,0) \leftarrow \ true \ F(1,1) \leftarrow \ true \ F(x,y_1+y_2) \leftarrow \ F(x-1,y_1) \wedge F(x-2,y_2) \wedge x > 1 \ false \leftarrow \ F(x,y) \wedge y < 0 \end{aligned}$$

Deep Learning-Based Framework (Instances)



- ☐ Guide Constrained Horn Clauses (CHCs) solving (Paper I and II)
- Guide word equation solving (Paper III and IV)

$$XaY = YbX \wedge XabY = YbaX$$

Constrained Horn Clauses (CHCs)

■ A CHC is a formula in the format

$$L(\bar{t}) \leftarrow L_1(\bar{t}_1) \wedge \ldots \wedge L_n(\bar{t}_n) \wedge \varphi$$

where \bar{t}_i are terms,

 L, L_1, \ldots, L_n are relation symbols,

 $L_i(\bar{t}_i)$ is an atom,

 φ is a constraint in the background theory T.

A CHC system is satisfiable if there exists a interpretation such that every clause in the system evaluates to true.

CHCs (Example)

Fibonacci function logic expression

$$egin{aligned} \operatorname{fib}(0) &= 0 \ \operatorname{fib}(1) &= 1 \ \operatorname{fib}(n) &= \operatorname{fib}(n-1) + \operatorname{fib}(n-2), \ \operatorname{for} \ n > 1 \end{aligned}$$

Its CHC encoding:

$$egin{aligned} F(0,0) \leftarrow \ true \ F(1,1) \leftarrow \ true \ F(x,y_1+y_2) \leftarrow \ F(x-1,y_1) \wedge F(x-2,y_2) \wedge x > 1 \end{aligned}$$

CHCs (Example)

An assertions such as "the Fibonacci function does not return negative numbers" can be encoded as:

$$false \leftarrow F(x,y) \land y < 0$$

Where false is a predicate representing an assertion violation.

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Satisfiability of CHCs (Example)

CHC encoding of Fibonacci function with an assertion:

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Consider a model:

$$F(x,y)\equiv x\geq 0 \land y\geq 0$$

Replacing F with this formula will make all clauses valid which means the system is satisfiable (SAT).

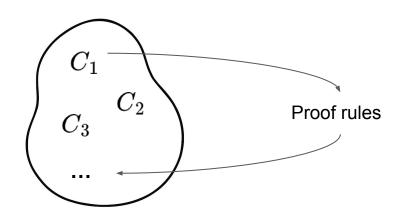
CHC Solver

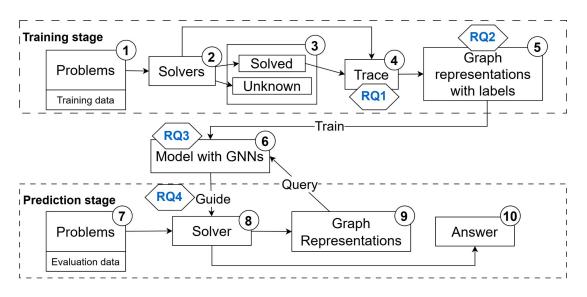
- Algorithms implemented in CHC solvers
 - Counterexample-Guided Abstraction Refinement (CEGAR).
 - Symbolic execution.

CHC Solver

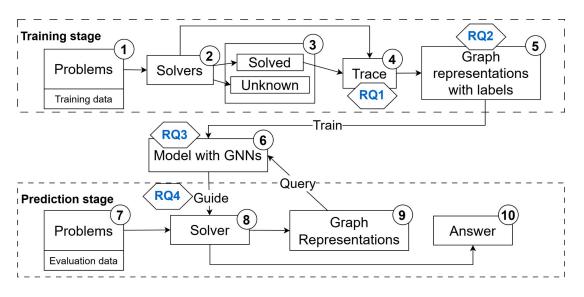
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Learning to tank clauses before solving



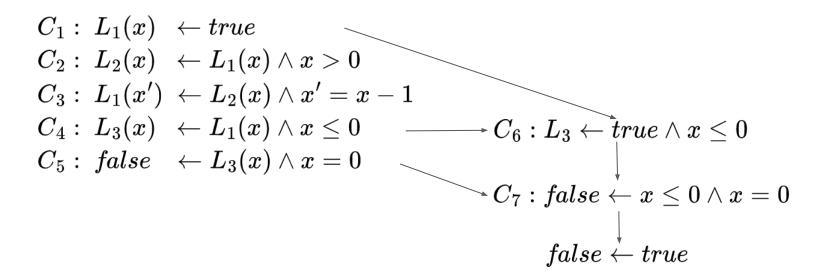


- RQI:Train task.
- RQ2: Graph representation.
- RQ3: GNN models.
- RQ4: Integrating methods.



- Answer to RQI (train task):
 - Learn from Minimal Unsatisfiable Subset (MUS).

Example of MUS



Example of MUS

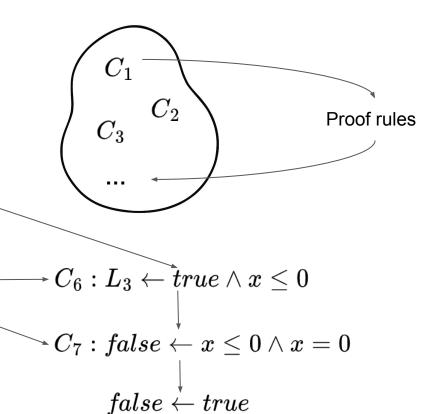
$$C_1: L_1(x) \leftarrow true$$

$$C_2: \ L_2(x) \ \leftarrow L_1(x) \land x > 0$$

$$C_3:\ L_1(x')\ \leftarrow L_2(x) \wedge x' = x-1$$

$$C_4:\ L_3(x)\ \leftarrow L_1(x) \wedge x \leq 0$$

$$C_5: \ false \ \leftarrow L_3(x) \wedge x = 0$$



Example of MUS

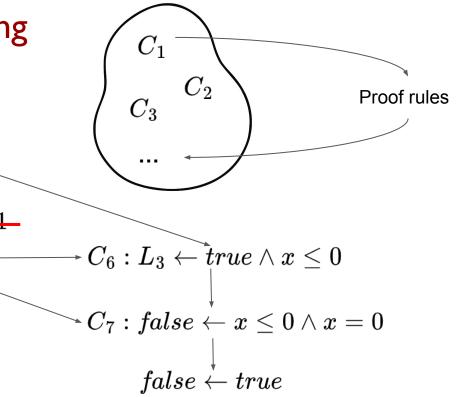
$$C_1: L_1(x) \leftarrow true \ C_2: L_2(x) \leftarrow L_1(x) \wedge x > 0$$

$$-C_3: L_1(x') \leftarrow L_2(x) \wedge x' - x - 1$$

$$C_4:\ L_3(x)\ \leftarrow L_1(x) \wedge x \leq 0$$

$$C_5: \ false \ \leftarrow L_3(x) \wedge x = 0$$

$$\{C1, C4, C5\}$$
 is a MUS.

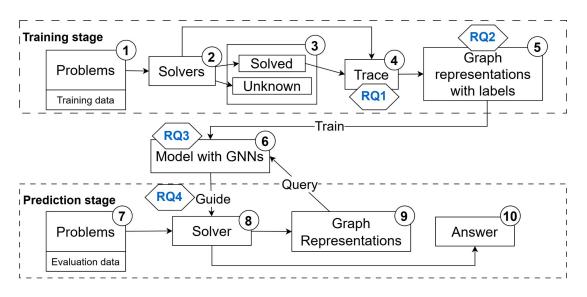


Example of MUS

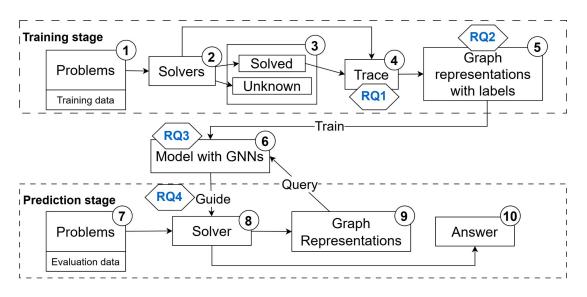
$$C_1: L_1(x) \leftarrow true$$
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Let ϕ be a set of formulas such that ϕ is unsatisfiable. A subset $\phi' \subseteq \phi$ is called a MUS if ϕ' is unsatisfiable and for all proper subset $\phi'' \subseteq \phi'$, ϕ'' is satisfiable.

Proof rules



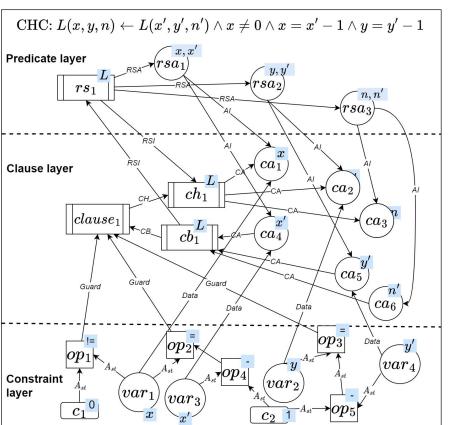
- Answer to RQI (train task):
 - Belongs to MUS or not.



- Answer to RQ2 (graph representation):
 - Syntactic structure.
 - Control-flow and data-flow.

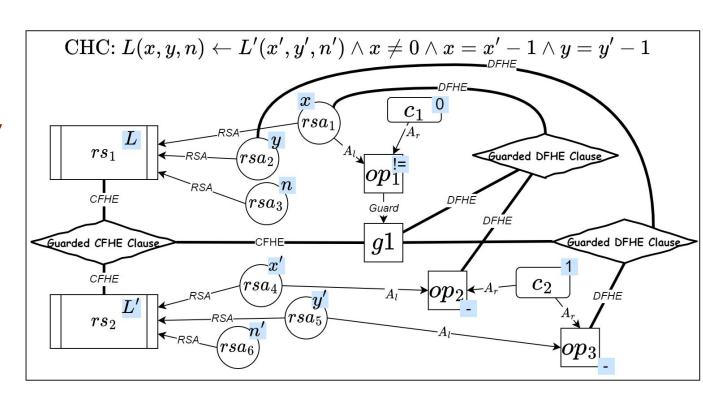
Rank CHCs to Guide the Solving (Graph Representation)

Constraint graph (CG)

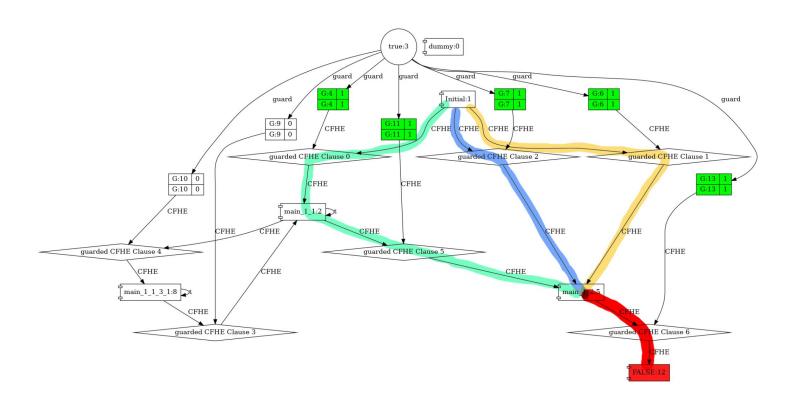


Rank CHCs to Guide the Solving (Graph Representation)

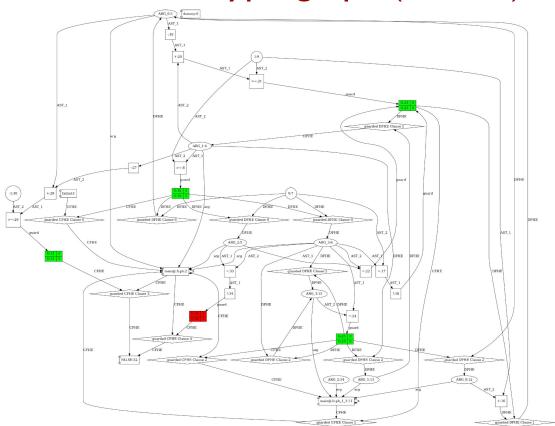
Control- and data- flow hypergraph (CDHG)

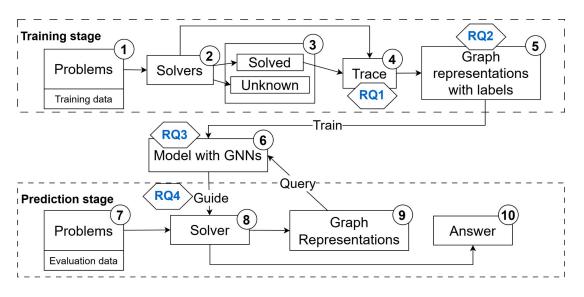


Control- and Data-Flow Hypergraph (CDHG)



Control- and Data-Flow Hypergraph (CDHG)





- Answer to RQ3 (GNN model):
 - Relational Hyper-Graph Neural Network (R-HyGNN).

Relational Hyper-Graph Neural Network (R-HyGNN)

The updating rule for node representation in time step t:

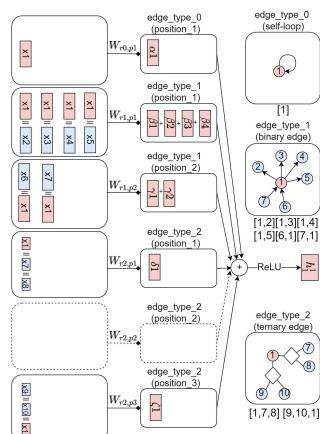
$$h_v^t = ext{ReLU}(\sum_{r \in R} \sum_{p \in P_r} \sum_{e \in E_v^{r,p}} W_{r,p}^t) \cdot || \{h_u^{t-1} | u \in e\},$$

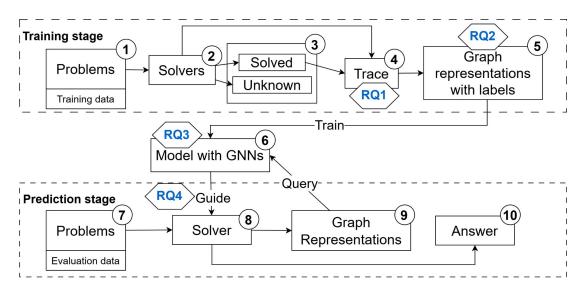
where $||\{\cdot\}|$ denotes concatenation of elements in a set, R is the set of edge types,

 P_r is the set of node position under edge type r,

 $W_{r,n}^t$ denotes learnable parameters,

 $E_{v}^{r,p}$ is the set of hyperedges.





- Answer to RQ4 (integrating methods):
 - Use prediction alone.
 - Combine with existing heuristics.
 - Combine with random clause selection.

Experimental Results (Improved Percentage)

- Evaluated in CHC solver Eldarica [6].
- Using CHC-COMP dataset.

Benchmark	MUS data set (best count)	Best ranking function (improvement in %)						
Algorithm		Number of Solved Problems			Average Time			
		Total	SAT	UNSAT	All	Common	SAT	UNSAT
5	Union	R-Plus	R-Plus	R-Minus	R-Plus	S-Plus	S-Minus	Rank
Linear	(0)	(1.4%)	(2.4%)	(1.0%)	(1.3%)	(19.1%)	(46.5%)	(31.1%)
CEGAR	Single	Rank	R-Plus	Rank	R-Plus	S-Plus	R-Minus	Rank
	(3)	(3.6%)	(4.0%)	(8.2%)	(1.9%)	(26.6%)	(57.9%)	(36.3%)
	Intersection	R-Plus	S-Plus	R-Plus	R-Plus	S-Plus	R-Minus	S-Plus
	(4)	(4.1%)	(0.8%)	(9.3%)	(3.1%)	(27.6%)	(45.0%)	(0.0%)
	Union	Two-Q	S-Plus*	Random	Two-Q	R-Minus	R-Minus	S-Plus
Linear	(4)	(1.0%)	(0.0%)	(2.0%)	(0.9%)	(12.7%)	(30.2%)	(26.5%)
SymEx	Single	S-Minus*	S-Plus*	Random	Random	S-Plus	Random	S-Plus
	(3)	(0.5%)	(0.0%)	(2.0%)	(0.8%)	(12.9%)	(28.4%)	(17.6%)
	Intersection	S-Plus*	S-Plus*	S-Plus*	S-Plus	Score	Random	R-Plus
	(5)	(1.0%)	(0.0%)	(2.0%)	(1.3%)	(9.5%)	(28.4%)	(35.8%)

Word Equations (example)

A word equation:

$$XabY = YbaX$$

where a and b are letters,

X,Y, and Z are variables ranging over strings of these letters.

Satisfiability of a Word Equation

A word equation:

$$XabY = YbaX$$

where a and b are letters, X,Y, and Z are variables ranging over strings of these letters.

If there exists a substitution of the variables with strings over the letters that makes the equation hold, then the equation is satisfiable.

Satisfiability of a Word Equation

A word equation:

$$XabY = YbaX$$

where a and b are letters, X,Y, and Z are variables ranging over strings of these letters.

If there exists a substitution of the variables with strings over the letters that makes the equation hold, then the equation is satisfiable.

SAT or UNSAT?

Satisfiability of a Word Equation

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where a and b are letters, X,Y, and Z are variables ranging over strings of these letters.

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$$X = b, Y = \epsilon$$

$$bab = bab$$

Word Equations

- Important in modeling string constraints in verification tasks.
 - ☐ E.g., Validate user inputs, ensuring correct string manipulations.

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Word Equations

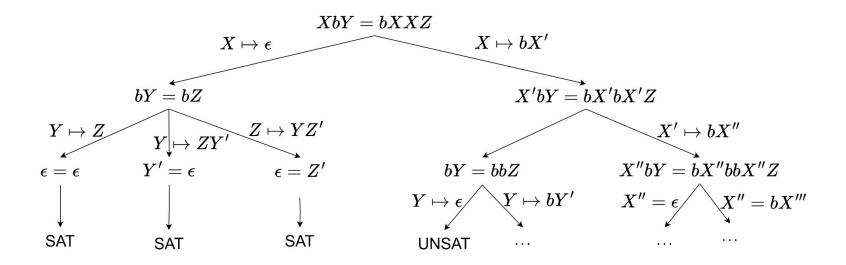
- Important in modeling string constraints in verification tasks.
 - ☐ E.g., Validate user inputs, ensuring correct string manipulations.
- Difficult to solve.
 - Decision procedures such as [6] and [7] have no implementation.
 - Practical algorithms are incomplete.

- [6] Makanin, G.S.: The problem of solvability of equations in a free semigroup. 1977
- [7] Plandowski, W.: Satisfiability of word equations with constants is in pspace. 1999.

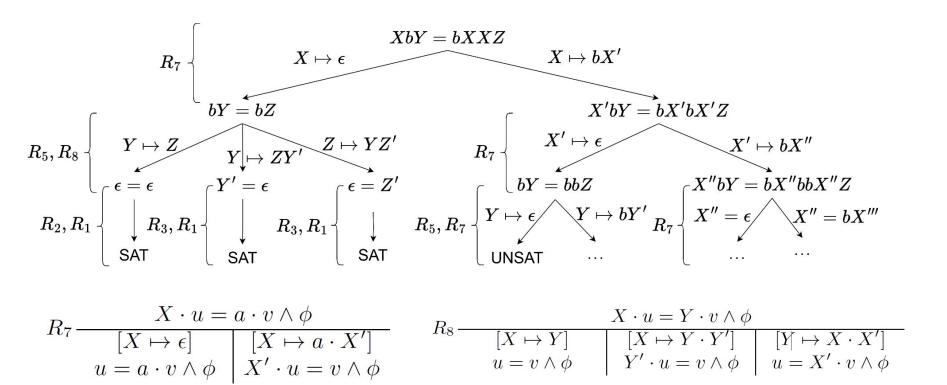
Solving a Word Equation System

- Split Algorithm
 - □ Based on Levi's lemma (Nielsen transformation in group theory).

Split Algorithm (Branch Process)



Split Algorithm (Branch Process)



Calculus

$$R_1 \frac{true}{\text{SAT}}$$
 $R_2 \frac{\epsilon = \epsilon \wedge \phi}{\phi}$ $R_3 \frac{X = \epsilon \wedge \phi}{[X \mapsto \epsilon]}$ $R_4 \frac{a \cdot u = \epsilon \wedge \phi}{\text{UNSAT}}$

 $R_5 \frac{a \cdot u = a \cdot v \wedge \phi}{u = v \wedge \phi}$

$$\phi$$
 with $X \in \Gamma$ and $a \in \Sigma$.

(a) Simplification rules

 $R_6 \frac{a \cdot u = b \cdot v \wedge \phi}{\text{UNSAT}}$

 $R_7 \frac{X \cdot u = a \cdot v \wedge \phi}{[X \mapsto \epsilon] \quad [X \mapsto a \cdot X']}$

with a, b two different letters from Σ . (b) Letter-letter rules

- $u = a \cdot v \wedge \phi$ $X' \cdot u = v \wedge \phi$ with X' a fresh element of Γ .
 - (c) Variable-letter rules

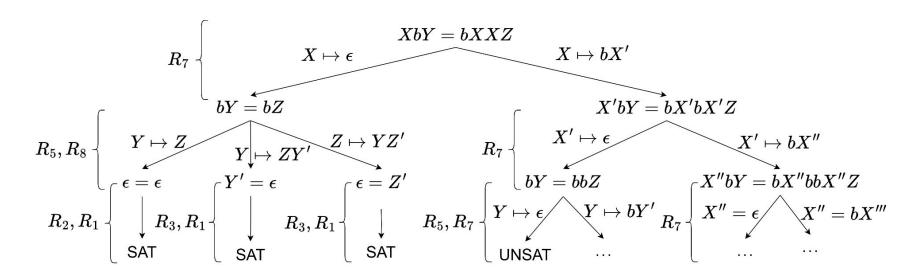
- $R_8 \frac{X \cdot u = Y \cdot v \wedge \phi}{[X \mapsto Y] \quad [X \mapsto Y \cdot Y'] \quad [Y \mapsto X \cdot X']}$
 - with $X \neq Y$ and X', Y' fresh elements of Γ .

$$R_9 \frac{X \cdot u = X \cdot v \wedge \phi}{u = v \wedge \phi}$$

(d) Variable-variable rules

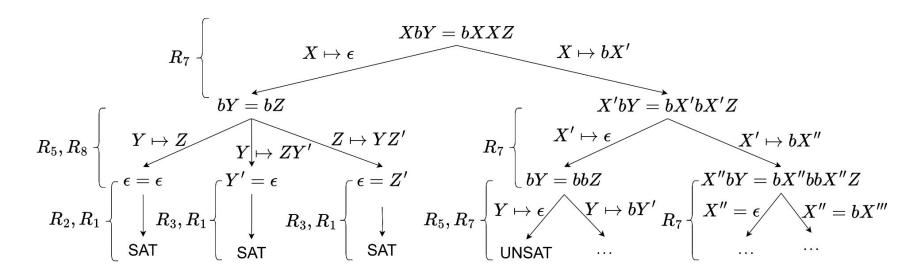
63

Split Algorithm (Branch Process)



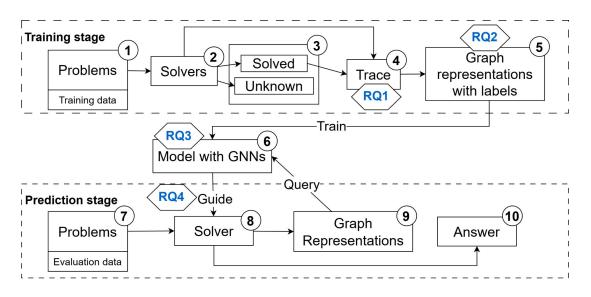
- If there is one branch SAT, then the word equation is SAT.
- If all branches are UNSAT, then the word equation is UNSAT.

Split Algorithm (Branch Process)



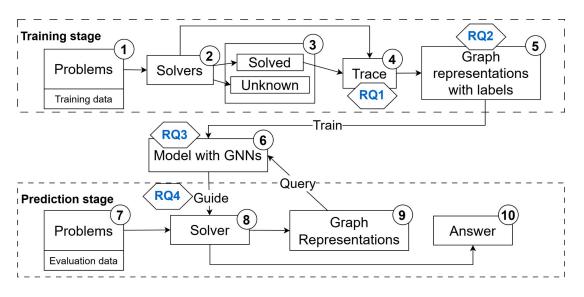
Branching significantly affects the performance.

Select Branches to Guide the Solving

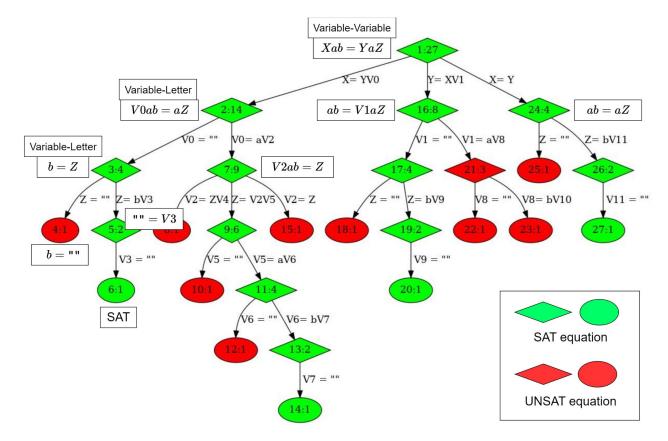


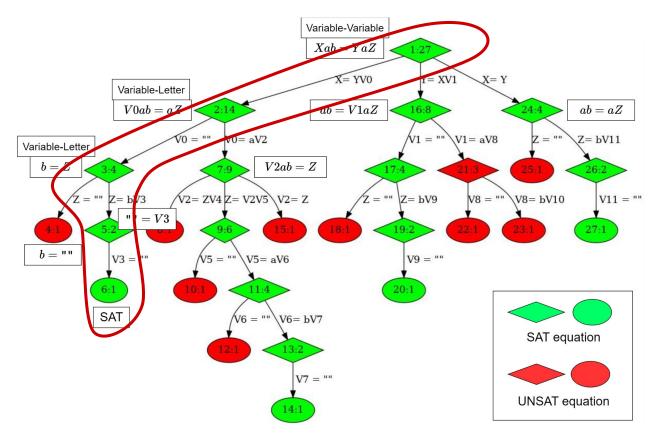
- RQI:Train task.
- RQ2: Graph representation.
- RQ3: GNN models.
- RQ4: Integrating methods.

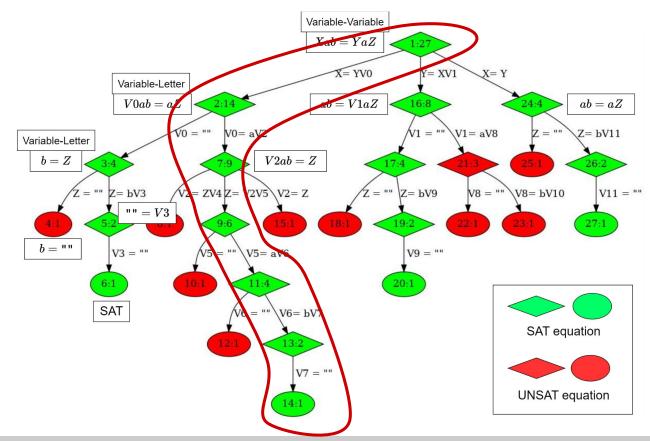
Select Branches to Guide the Solving

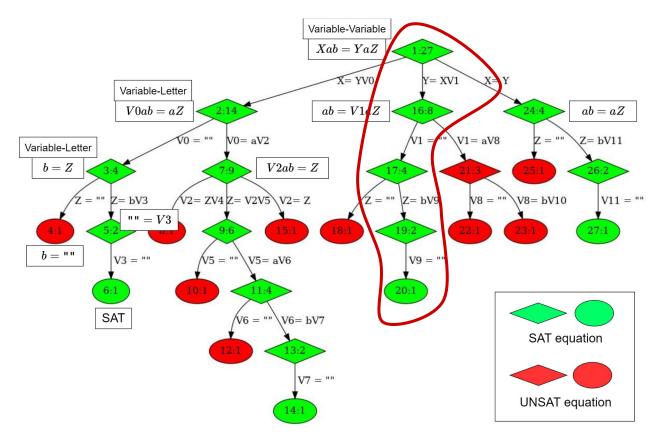


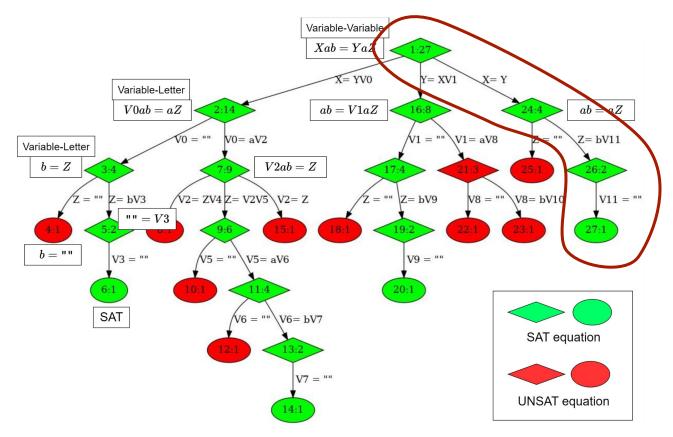
- Answer to RQI (train task):
 - Learn from shortest path to SAT.



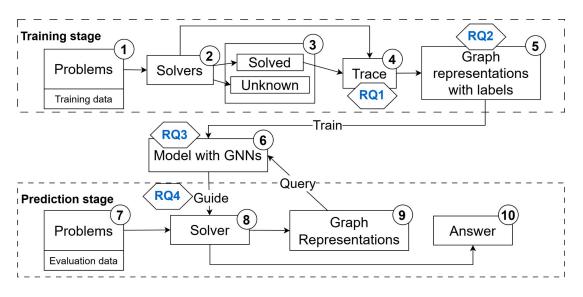








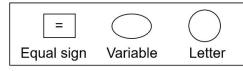
Select Branches to Guide the Solving

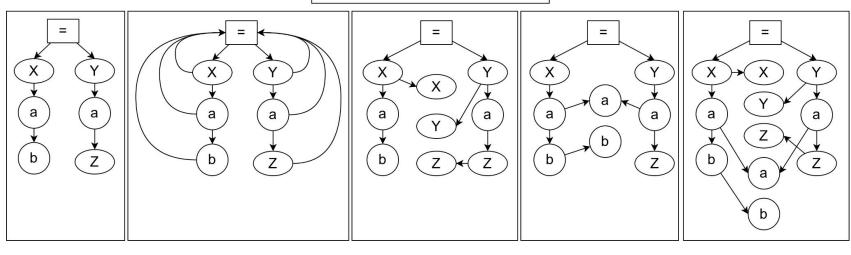


Answer to RQ2 (graph representation).

Select Branches to Guide the Solving (Graph Representation)

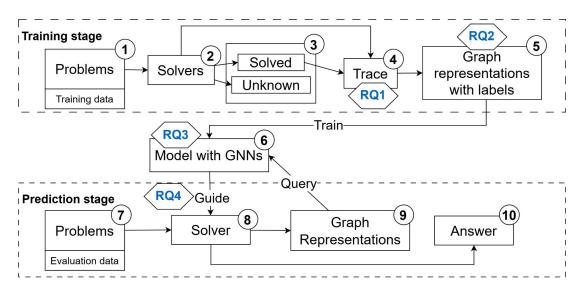
$$Xab = YaZ$$





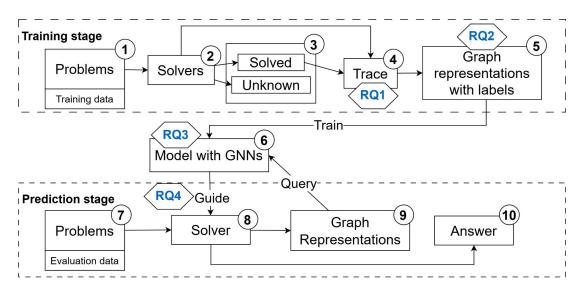
Graph 1 Graph 2 Graph 3 Graph 4 Graph 5

Select Branches to Guide the Solving



- Answer to RQ3 (GNN model):
 - Graph Convolutional Network (GCN).
 - Graph Attention Network (GAT).
 - Graph Isomorphism Network (GIN).

Select Branches to Guide the Solving



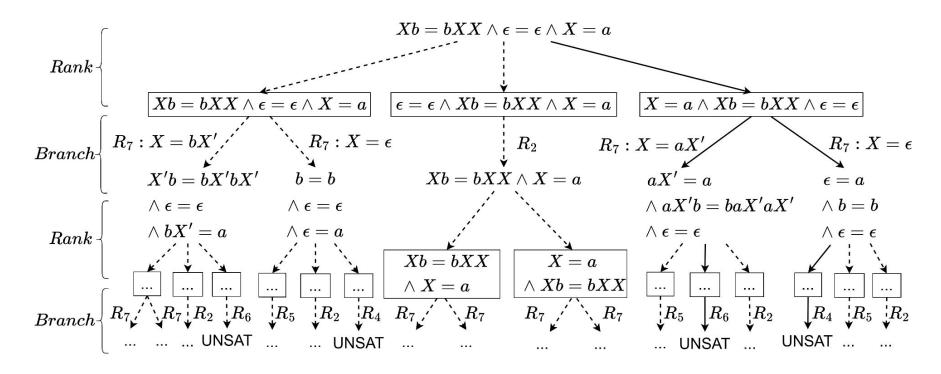
- Answer to RQ4 (integrating methods):
 - Use prediction alone.
 - Combine with random branch selection.

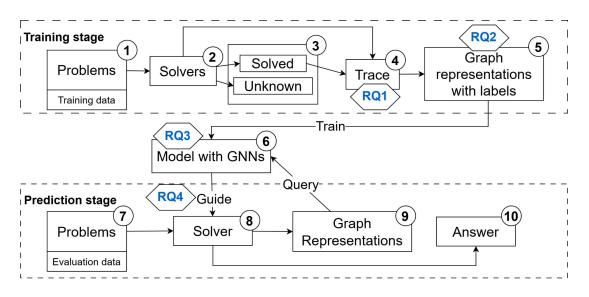
Experimental Results

■ Evaluated in our word equation solver, Z3, cvc5, etc.

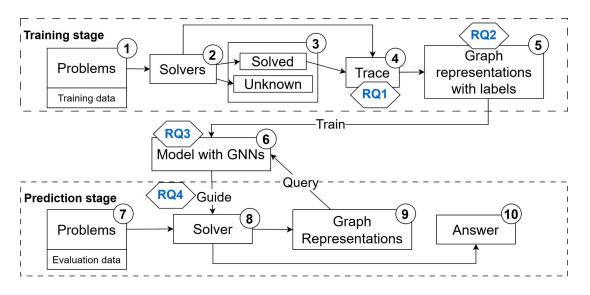
Bench	Solver	Number of solved problems					Average solving time (split number)				
		SAT	UNS	UNI	CS	CU	SAT	UNS	CS	CU	
2 (1000 in total)	Fixed	33	0	10		0	13.2 (1115.2)	- (-)	4.8 (7)	- (-)	
	Random	41	0	6			11.7 (3879.5)	- (-)	4.2 (60)	- (-)	
	GNN	71	0	27	1		46.0 (1813.5)	- (-)	5.1 (5)	- (-)	
	cvc5	4	2	4			2.0 (-)	0.1 (-)	0.1 (-)	- (-)	
	Ostrich	14	43	44			40.7 (-)	31.8 (-)	2.5 (-)	- (-)	
	Woorpje	23	0	2			38.3 (-)	- (-)	0.1 (-)	- (-)	
	Z3	6	0	2			0.1 (-)	- (-)	4.2 (-)	- (-)	
	Z3-Noodler	19	0	0			45.8 (-)	- (-)	4.2 (-)	- (-)	

Solving Word Equation System



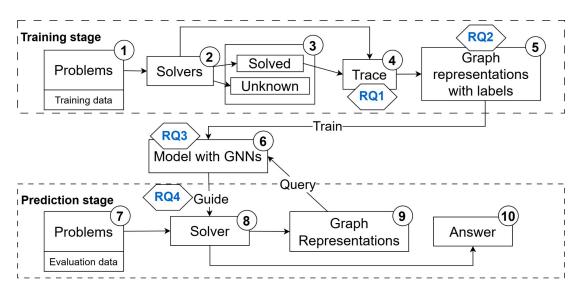


- RQI:Train task.
- RQ2: Graph representation.
- RQ3: GNN models.
- RQ4: Integrating methods.

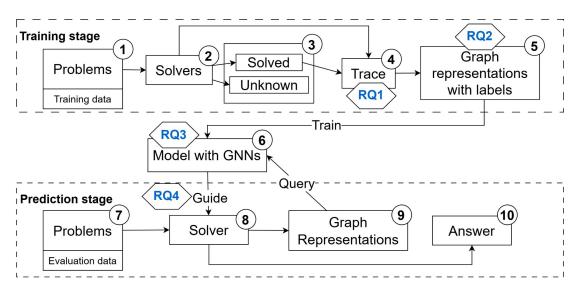


Answer to RQI (train task):

$$XaY = YbX \land XabY = YbaX$$



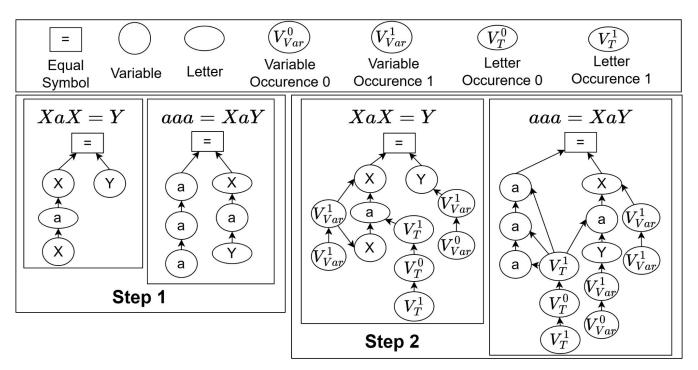
- Answer to RQI (train task):
 - Learn from MUSes given from other solvers.
 - Learn from shortest path from split algorithm.



- Answer to RQ2 (graph representation):
 - □ Global information.

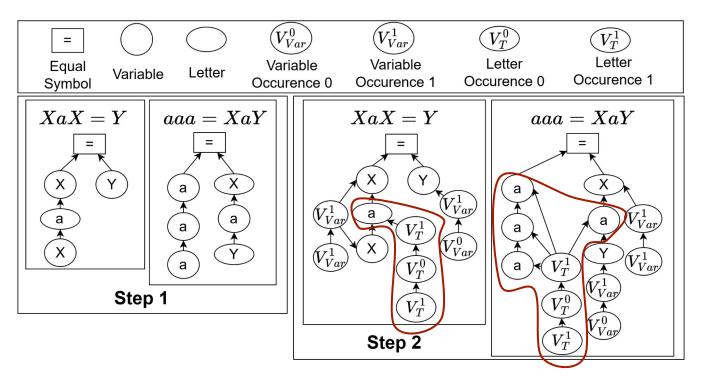
Rank Word Equations (Graph Representation)

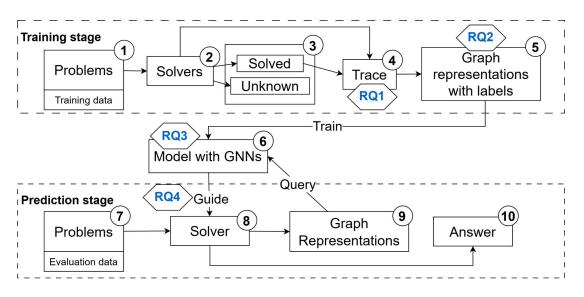
$$XaX = Y \wedge aaa = XaY$$



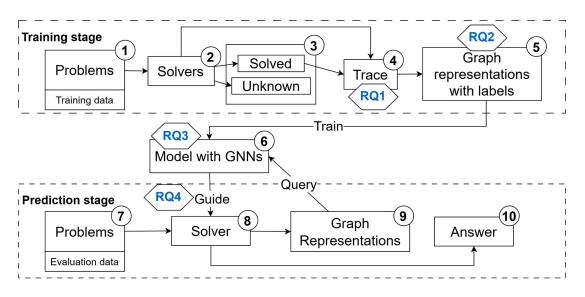
Rank Word Equations (Graph Representation)

$$XaX = Y \wedge aaa = XaY$$





- Answer to RQ3 (GNN model):
 - GCN.
 - □ GCN+GIN.
 - ☐ GNN filters.



- Answer to RQ4 (integrating method):
 - □ Different frequency (e.g., one-short).
 - Combine with random ranking strategy.
 - Combine with manually designed heuristics.

Experimental Results

■ Evaluated in our word equation solver, Z3, cvc5, etc.

Bench	Solver	Number of solved problems					Average solving time (split number)			
		SAT	UNSAT	UNI	CS	CU	SAT	UNSAT	CS	CU
A1	DragonLi	24	955	0	13	678	5.6 (244.8)	6.5 (1085.3)	5.0 (94.4)	5.7 (126.3)
	Random- DragonLi	22	944	0			5.6 (198.8)	6.3 (932.6)	5.6 (137.6)	5.7 (180.5)
	GNN- DragonLi	24	961	0			6.1 (164.7)	7.5 (1974.8)	6.1 (96.4)	6.3 (60.5)
	cvc5	24	952	1			0.5	0.6	0.1	0.3
	Z3	17	960	0			8.7	0.4	1.1	0.1
	Z3-Noodler	22	939	2			5.7	0.3	4.8	0.1
	Ostrich	17	931	0			15.0	5.5	8.0	4.7
	Woorpje	23	744	0			3.0	12.5	0.1	12.2

- RQI:What are good encodings of symbolic decision processes as training tasks?
 - Encode the problems to classification task.
 - Collect train data from multiple sources.

- □ RQI:What are good encodings of symbolic decision processes as training tasks?
 - Encode the problem to classification task.
 - ☐ Collect train data from multiple sources.
- RQ2:What is the most effective format for representing formulas in deep learning?
 - ☐ The graph representation must include all syntactic elements.
 - Use compact graph encoding (e.g., merge identical nodes).

RQI:What are good encodings of symbolic decision processes as training tasks? Encode the problem to classification task. Collect train data from multiple sources. □ RQ2:What is the most effective format for representing formulas in deep learning? The graph representation must include all syntactic elements. Use compact graph encoding (e.g., merge identical nodes). RQ3: Which deep learning technique is best suited for feature extraction from formulas? GCN serves as baseline. GAT, GIN, R-HyGNN, etc.

Ч	RQ1:	What are good encodings of symbolic decision processes as training tasks?
		Encode the problem to classification task.
		Collect train data from multiple sources.
	RQ2:\	What is the most effective format for representing formulas in deep learning?
		The graph representation must include all syntactic elements.
		Use compact graph encoding (e.g., merge identical nodes).
	RQ3:\	Which deep learning technique is best suited for feature extraction from formulas?
		GCN serves as baseline.
		GAT, GIN, R-HyGNN, etc.
	RQ4:\	What are the methods for integrating the trained model into algorithms?
		Cache embeddings.
		Combine with predefined heuristics did not yield the best performance.

Conclusions

- A deep learning-based framework for decision problems in symbolic methods.
- Two instances of the framework.
 - CHC solver.
 - Word equation solver.

Conclusions

- A deep learning-based framework for decision problems in symbolic methods.
- Two instances of the framework
 - CHC solver.
 - Word equation solver.
- Our framework can easily adapt to other decision problems in symbolic methods.
- GNN is still the best option for extracting structural information in symbolic expression.

Future Directions

- Training tasks
 - Reinforcement learning.
 - Sequential model.
 - Generative model.
- Extend to new problem domains.
 - Regular expression in word equations.
 - New theories in CHCs.
- Simultaneously guide multiple decision processes.



Thank you for listening

Word Equation System

A word equation:

$$XabY = YbaX$$

where a and b are letters, X,Y, and Z are variables ranging over strings of these letters.

A word equation system (conjunctive word equations):

$$\phi = e_1 \wedge \ldots \wedge e_n$$

where e_i is a word equation.

$$XaY = YbX \wedge XabY = YbaX$$

Motivating Examples

Prove:

$$egin{aligned} &orall x.\, L_1(x) \leftarrow true \ \land \ &orall x.\, L_2(x) \leftarrow \land x > 0 \ \land \ &orall x, x'.\, L_1(x') \leftarrow L_2(x) \land x' = x - 1 \ \land \ &orall x.\, L_3(x) \leftarrow L_1(x) \land x \leq 0 \ \land \ &orall x.\, false \leftarrow L_3(x) \land x = 0 \end{aligned}$$

For Simplicity:

$$egin{array}{lll} C1: \ L_1(x) &\leftarrow true \ C2: \ L_2(x) &\leftarrow L_1(x) \wedge x > 0 \ C3: \ L_1(x') &\leftarrow L_2(x) \wedge x' = x - 1 \ C4: \ L_3(x) &\leftarrow L_1(x) \wedge x \leq 0 \ C5: \ false &\leftarrow L_3(x) \wedge x = 0 \end{array}$$

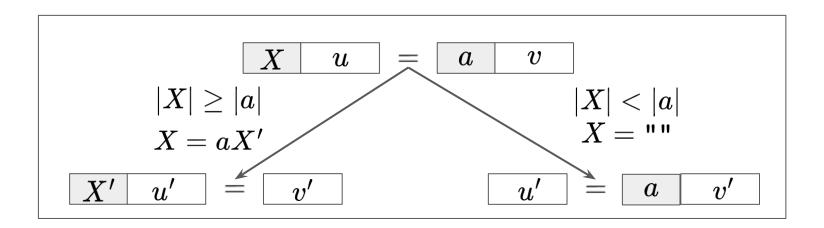
Motivating Examples (Program Verification)

Prove:

$$egin{aligned} &orall x.\, L_1(x) \leftarrow true \ \land \ &orall x.\, L_2(x) \leftarrow \land x > 0 \ \land \ &orall x,\, x'.\, L_1(x') \leftarrow L_2(x) \land x' = x - 1 \ \land \ &orall x.\, L_3(x) \leftarrow L_1(x) \land x \leq 0 \ \land \ &orall x.\, false \leftarrow L_3(x) \land x = 0 \end{aligned}$$

Solving a Word Equation System

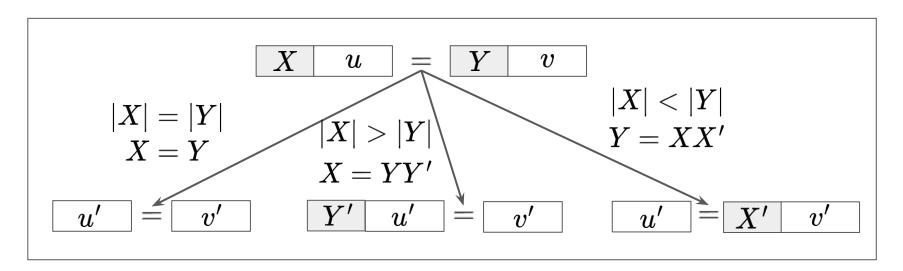
- Split Algorithm (branching example)
 - First terms are variable and terminal



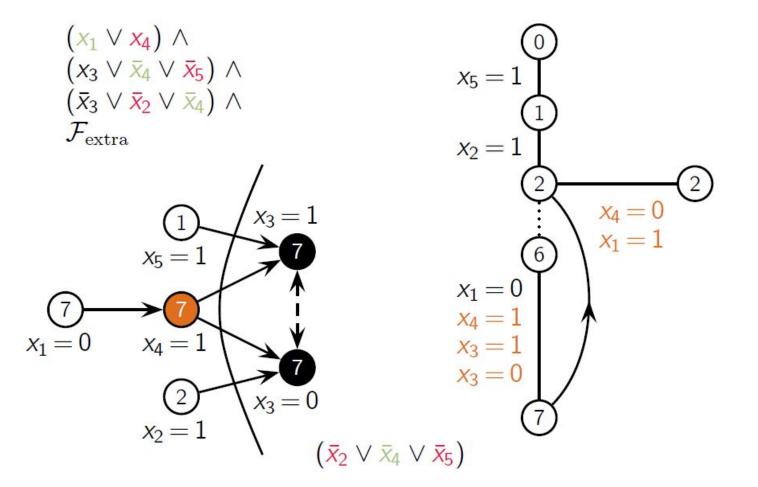
u, v, u', v' are terms, a is a letter, and X, X' are variables

Solving a Word Equation System

- Split Algorithm (branching example)
 - First terms are variables

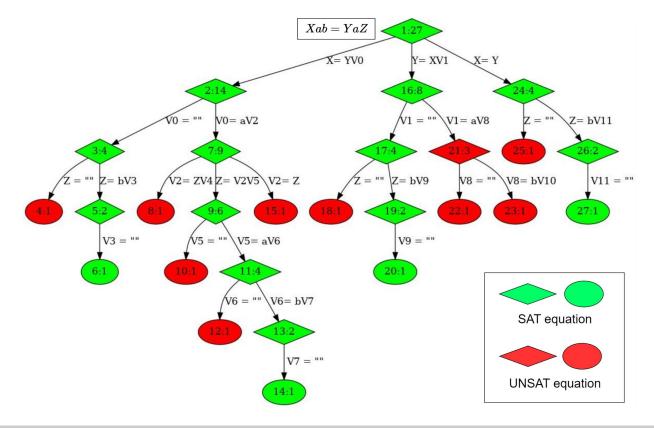


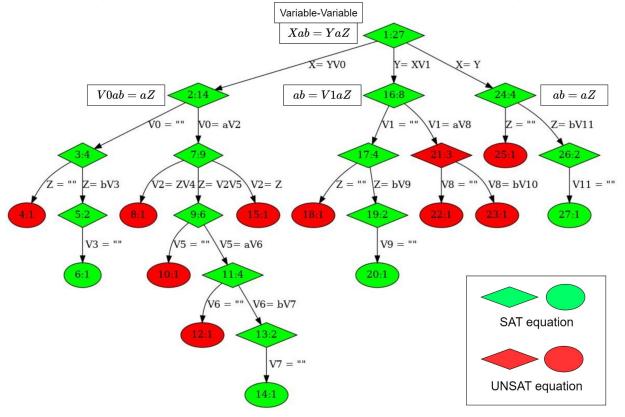
u, v, u', v' are terms, and X, Y, X', Y' are variables

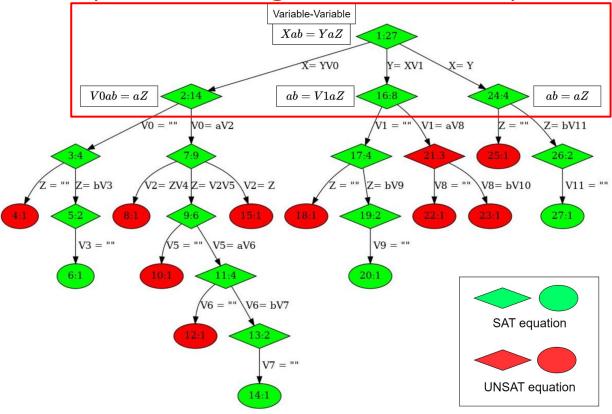


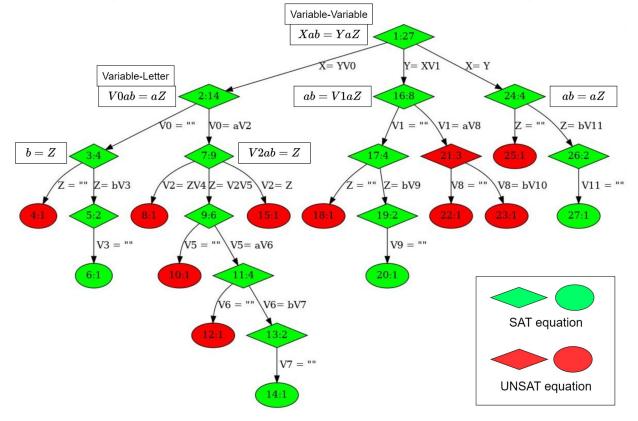
Symbolic Methods

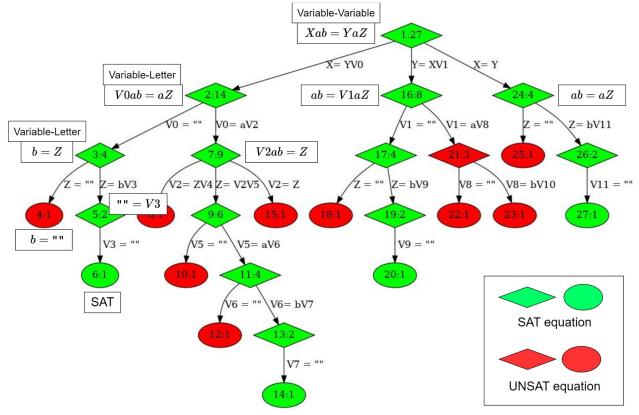
- Approaches
 - Automatic Theorem Provers (ATPs)
 - An example to show why this is interesting
 - Boolean Satisfiability (SAT)/Satisfiability Modulo Theories (SMT) solvers
 - Constrained Horn Clause (CHC) solvers
- Challenges
 - Complex representation
 - Theory handling
 - Scalability



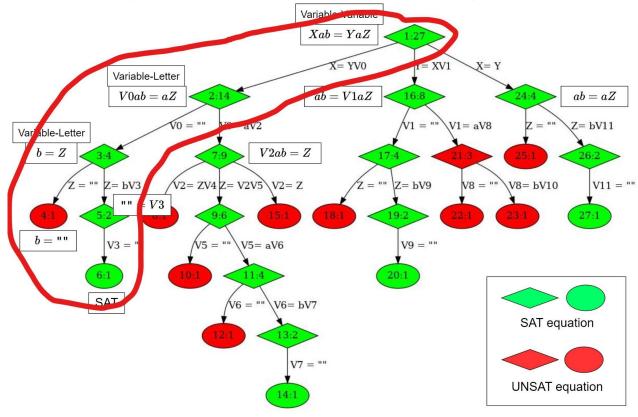




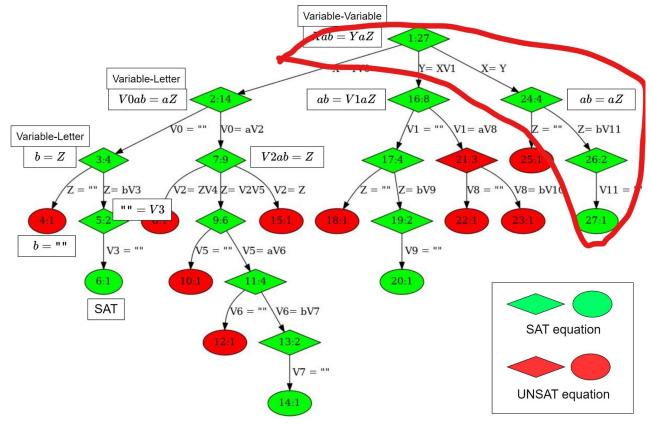




Split Algorithm (A Path Leading to SAT)

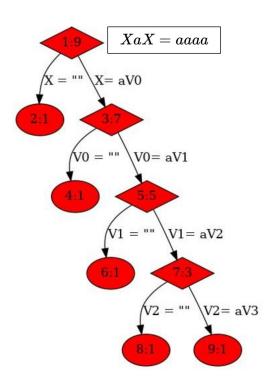


Split Algorithm (A Shortest Path Leading to SAT)

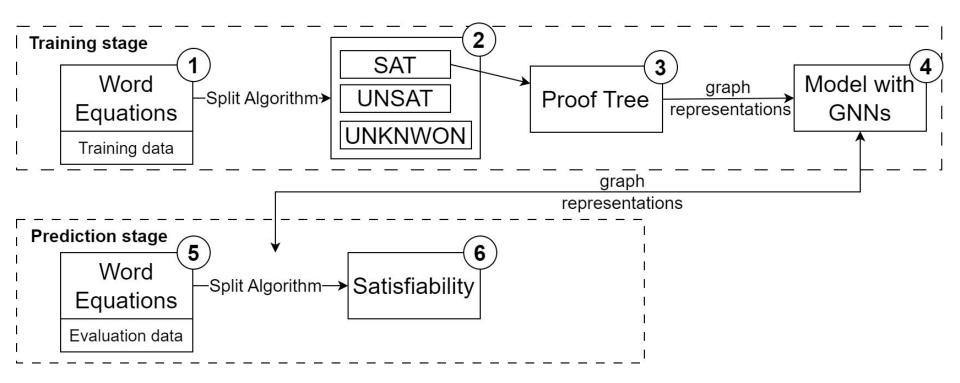


Split Algorithm (Proof Tree for UNSAT Problem)

 Need explore all paths to conclude UNSAT



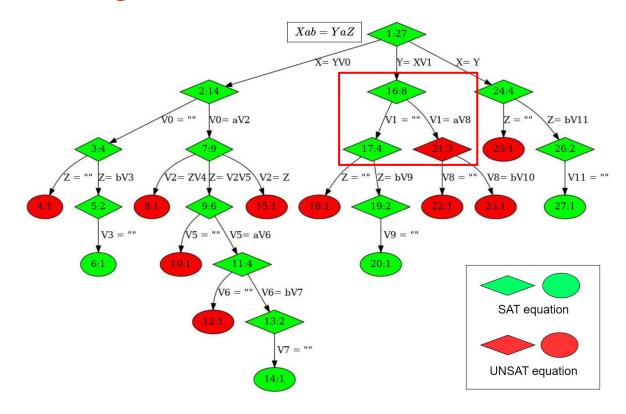
Working Pipeline (Data-driven Based Heuristic)



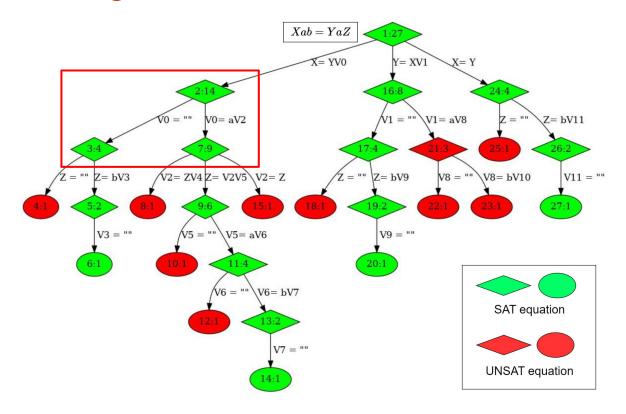
Graph Neural Networks (GNNs)

- A set of fully connected neural networks
- Take graph as input, output node and graph representations
- Can capture the structural features of graph

Label the Training Data at a Branch Point

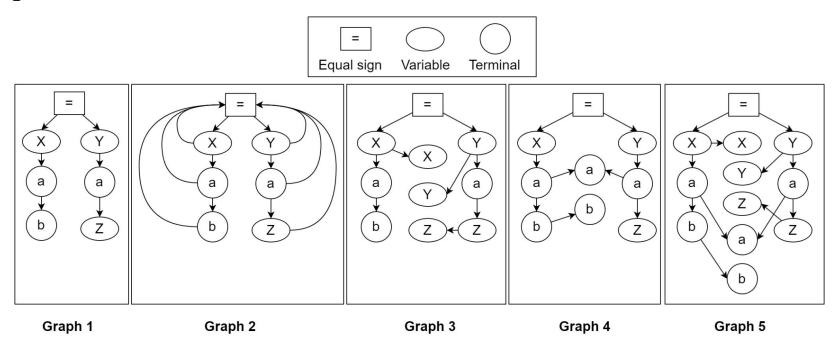


Label the Training Data at a Branch Point

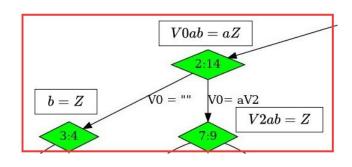


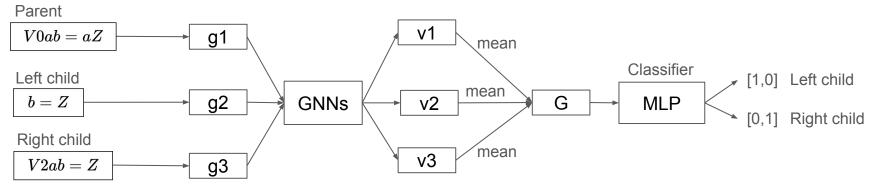
Graph Representations of Word Equation

Equation: Xab = YaZ



Model Structure





1. Randomly generated SAT word equation

aksCwqeuafhkajshfweeta = aksfjdfbabDeuafhkajshBCDta

- 1. Randomly generated SAT word equation
- 2. Word equation with a particular pattern [5]

$$X_naX_nbX_{n-1}bX_{n-2}\cdots bX_1=aX_nX_{n-1}X_{n-1}bX_{n-2}b\cdots bX_1X_1baa$$

[5] Day, J.D., Ehlers, T., Kulczynski, M., Manea, F., Nowotka, D., Poulsen, D.B.: On solving word equations using SAT. Reachability Problems, 93–106. (2019)

- 1. One randomly generated SAT word equation
- 2. One word equation with a particular pattern
- 3. Conjunctive word equations of Benchmark 1

$$egin{aligned} kSY &= WHXGk \ & \land vZX = vtntssemtvnetm \ & \land yffyyfFWVff = yffyyfyXLGZfU \ & \land NRQxgGI = xxgggFJTK \end{aligned}$$

- 1. One randomly generated SAT word equation
- 2. One word equation with a particular pattern
- 3. Conjunctive word equations of Benchmark 1
- QF_S, QF_SLIA, and QF_SNLIA tracks of SMT-LIB without length constraints, regular expressions, and Boolean operators

$$AabcdBCefDEghF = BciabADGCFHE \ abAaBabcdCdcDE = AbeFCcdbaBDcdGfg$$

- 1. One randomly generated SAT word equation
- 2. One word equation with a particular pattern
- 3. Conjunctive word equations of Benchmark 1
- 4. QF_S, QF_SLIA, and QF_SNLIA tracks of SMT-LIB without length constraints, regular expressions, and Boolean operators

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Ben To			rk 1 000			chmark al: 2100		500		mark 3 : 41000		100001110011		mar : 23	2.02 (YAS
20	00		Eval	2	20000			40000			Eval	1	855		Eval
√	×	∞	Evai	√	×	∞	Eval	√	×	∞	Evai	√	×	∞	Livai
1997	0	3	1000	1293	0	18707	1000	1449	1137	37414	1000	1673	16	166	455

		Nι	ımber	of so	olvec	l		Aver	age	-
Bench	Solver		pro	blems	S		solvi	ng time (s	plit num	ber)
		SAT	UNS	UNI	CS	CU	SAT	UNS	CS	CU
	Fixed	999	-	0			4.1 (182.0)	- (-)	4.0 (169.0)	- (-)
1	Random	996	_	0			4.2 (349.6)	- (-)	4.1 (269.8)	- (-)
(1000 SAT)	GNN	995	1	0	777	0	7.6 (215.7)	- (-)	7.0 (162.3)	- (-)
	cvc5	1000	-	0			0.1 (-)	- (-)	0.1 (-)	- (-)
	Ostrich	918	-	0			20.4 (-)	- (-)	19.6 (-)	- (-)
	Woorpje	967	-	0			1.6 (-)	- (-)	0.5 (-)	- (-)
	Z3	902	_	0			3.4 (-)	- (-)	2.4 (-)	- (-)
	Z3-Noodler	935	-	0			1.9 (-)	- (-)	1.1 (-)	- (-)

		Nι	ımber	of so	olvec	ł		Aver	age	
Bench	Solver		pro	blems	S		solvi	ng time (s	plit num	ber)
		SAT	UNS	UNI	CS	CU	SAT	UNS	CS	CU
	Fixed	999	-	0			4.1 (182.0)	- (-)	4.0 (169.0)	- (-)
1	Random	996	-	0			4.2 (349.6)	- (-)	4.1 (269.8)	- (-)
(1000 SAT)	GNN	995	1	0	777	0	7.6 (215.7)	- (-)	7.0 (162.3)	- (-)
	cvc5	1000	-	0			0.1 (-)	- (-)	0.1 (-)	- (-)
	Ostrich	918	-	0			20.4 (-)	- (-)	19.6 (-)	- (-)
	Woorpje	967	-	0			1.6 (-)	- (-)	0.5 (-)	- (-)
	Z3	902	_	0			3.4 (-)	- (-)	2.4 (-)	- (-)
	Z3-Noodler	935	-	0			1.9 (-)	- (-)	1.1 (-)	- (-)

		Nι	ımber	of so	olvec	l		Aver	age	
Bench	Solver		pro	blems	S		solvi	ng time (s	plit num	ber)
		SAT	UNS	UNI	CS	CU	SAT	UNS	CS	CU
	Fixed	999	-	0			4.1 (182.0)	- (-)	4.0 (169.0)	- (-)
1	Random	996	-	0			4.2 (349.6)	- (-)	4.1 (269.8)	- (-)
(1000 SAT)	GNN	995	1	0	777	0	7.6 (215.7)	- (-)	7.0 (162.3)	- (-)
100	cvc5	1000	-	0			0.1 (-)	- (-)	0.1 (-)	- (-)
	Ostrich	918	-	0			20.4 (-)	- (-)	19.6 (-)	- (-)
	Woorpje	967	-	0			1.6 (-)	- (-)	0.5 (-)	- (-)
	Z3	902	_	0			3.4 (-)	- (-)	2.4 (-)	- (-)
	Z3-Noodler	935	-	0			1.9 (-)	- (-)	1.1 (-)	- (-)

7.		Nι	ımbeı	of so	olvec	ł		Aver	age	
Bench	Solver		pro	blems	S		solvi	ng time (s	plit num	ber)
		SAT	UNS	UNI	CS	CU	SAT	UNS	CS	CU
	Fixed	999	1	0			4.1 (182.0)	- (-)	4.0 (169.0)	- (-)
1	Random	996	-	0			4.2 (349.6)	- (-)	4.1 (269.8)	- (-)
(1000 SAT)	GNN	995	-	0	777	0	7.6 (215.7)	- (-)	7.0 (162.3)	- (-)
	cvc5	1000	·	0			0.1 (-)	- (-)	0.1 (-)	- (-)
	Ostrich	918	<u></u> -	0			20.4 (-)	- (-)	19.6 (-)	- (-)
	Woorpje	967	1-	0			1.6 (-)	- (-)	0.5 (-)	- (-)
	Z3	902	-	0			3.4 (-)	- (-)	2.4 (-)	- (-)
	Z3-Noodler	935	-	0			1.9 (-)	- (-)	1.1 (-)	- (-)

		Nι	ımber	of so	olvec	ł		Aver	age	
Bench	Solver		pro	blems	S		solvii	ng time (s	plit num	ber)
		SAT	UNS	UNI	CS	CU	SAT	UNS	CS	CU
	Fixed	33	0	10			13.2 (1115.2)	- (-)	4.8 (7)	- (-)
2 (1000	Random	41	0	6			11.7 (3879.5)	- (-)	4.2 (60)	- (-)
in total)	GNN	71	0	27	1	0	46.0 (1813.5)	- (-)	5.1 (5)	- (-)
total	cvc5	4	2	4			2.0 (-)	0.1 (-)	0.1 (-)	- (-)
	Ostrich	14	43	44			40.7 (-)	31.8 (-)	2.5 (-)	- (-)
	Woorpje	23	0	2			38.3 (-)	- (-)	0.1 (-)	- (-)
	Z3	6	0	2			0.1 (-)	- (-)	4.2 (-)	- (-)
	Z3-Noodler	19	0	0	,		45.8 (-)	- (-)	4.2 (-)	- (-)

<i>i</i>		Nι	ımber	of so	olvec	ł	Average						
Bench	Solver		pro	blems	S		solvi	ng time (s	plit num	ber)			
		SAT	UNS	UNI	CS	CU	SAT	UNS	CS	CU			
	Fixed	32	79	0			5.2	65.8	3.6	38.3			
	Tixed	02	10	U			(1946.2)	(4227.0)	(57.0)	(796.6)			
3	Random	32	79	0			9.5	65.0	3.8	38.5			
(1000)	Tandom	02	13	U			(3861.8)	(4227.0)	(61.7)	(796.6)			
in	GNN	32	65	0	23	50	214.3	1471.2	4.6	84.0			
total)	GIVIV	34	00	U			(1471.2)	(1471.2)	(63.7)	(796.6)			
total	cvc5	32	943	2			0.1 (-)	0.3 (-)	0.1 (-)	0.3 (-)			
	Ostrich	27	926	0			5.8 (-)	4.7 (-)	4.6 (-)	4.5 (-)			
	Woorpje	34	723	1			12.4 (-)	12.3 (-)	0.1 (-)	23.2 (-)			
	Z3	26	953	10			5.6 (-)	0.5 (-)	4.7 (-)	0.1 (-)			
	Z3-Noodler	28	926	0			22.7 (-)	0.3 (-)	8.9 (-)	0.1 (-)			

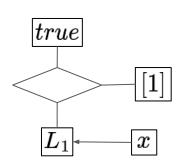
		Nι	ımber	of so	olvec	ł	Average						
Bench	Solver		prol	blems	S		solvi	ng time (s	plit num	ber)			
		SAT	UNS	UNI	CS	CU	SAT	UNS	CS	CU			
	Fixed	416	6	0	3		5.1 (105.5)	17.7 (17119.5)	5.1 (51.0)	5.0 (246)			
4	Random	415	6	0			4.9 (61.3)	17.9 (17119.5)	4.9 (38.1)	4.4 (246)			
(455 in	GNN	418	5	0	403	2	5.5 (118.3)	31.8 (5019.6)	5.3 (49.0)	8.8 (246)			
total)	cvc5	406	34	0			0.1 (-)	0.1 (-)	0.1 (-)	0.1 (-)			
	Ostrich	406	6	0			1.4 (-)	1.2 (-)	1.4 (-)	1.2 (-)			
	Woorpje	420	2	0			0.2 (-)	3.6 (-)	0.2 (-)	3.6 (-)			
	Z3	420	10	0			0.1 (-)	0.1 (-)	0.1 (-)	0.1 (-)			
	Z3-Noodler	420	35	1			0.1 (-)	0.1 (-)	0.1 (-)	0.1 (-)			

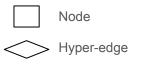
Extract train data

- Binary classification label
 - Union
 - Intersection
 - Single

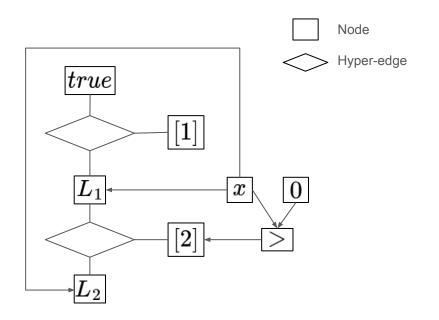
Label	Cl	lauses
1	[1] $L_1(x)$	$\leftarrow true$
0	$[2]\ L_2(x)$	$\leftarrow L_1(x) \land x > 0$
0	$[3] \ L_1(x')$	$\leftarrow L_2(x) \wedge x' = x - 1$
1	[4] $L_3(x)$	$\leftarrow L_1(x) \land x \leq 0$
1	[5] <i>false</i>	$\leftarrow L_3(x) \wedge x eq 0$

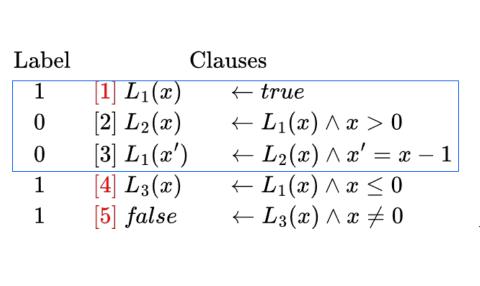
Label	Clauses	
1	$[1] L_1(x)$	$\leftarrow true$
0	$[2]\ L_2(x)$	$\leftarrow L_1(x) \land x > 0$
0	$[3] \ L_1(x')$	$\leftarrow L_2(x) \wedge x' = x - 1$
1	[4] $L_3(x)$	$\leftarrow L_1(x) \land x \leq 0$
1	$[5] \ false$	$\leftarrow L_3(x) \land x \neq 0$

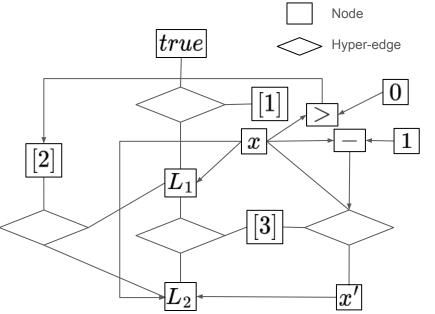


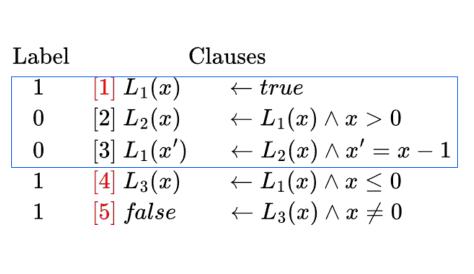


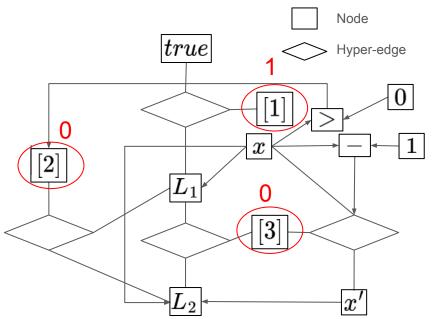
Label	Clauses		
1	$[1] L_1(x)$	$\leftarrow true$	
0	$[2]\ L_2(x)$	$\leftarrow L_1(x) \land x > 0$	
0	$[3] L_1(x')$	$\leftarrow L_2(x) \wedge x' = x - 1$	
1	$[4] \ L_3(x)$	$\leftarrow L_1(x) \land x \leq 0$	
1	[5] <i>false</i>	$\leftarrow L_3(x) \land x \neq 0$	

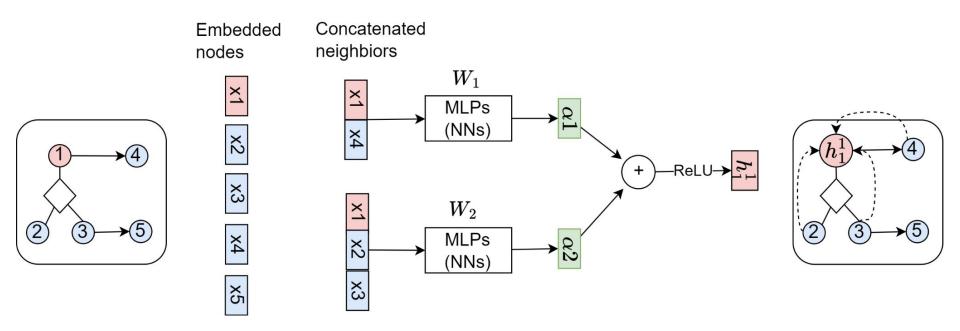


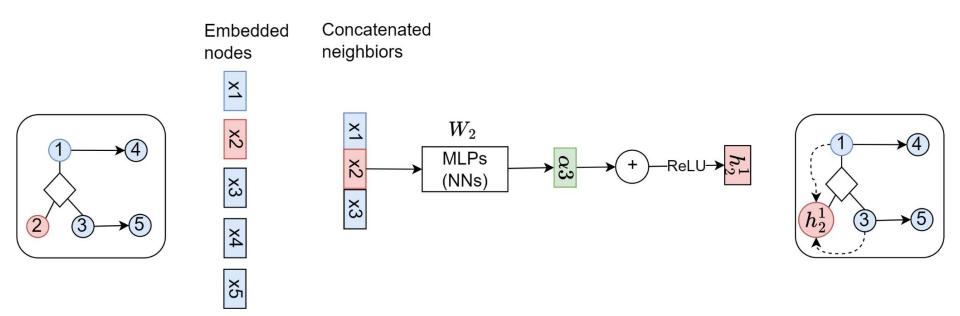


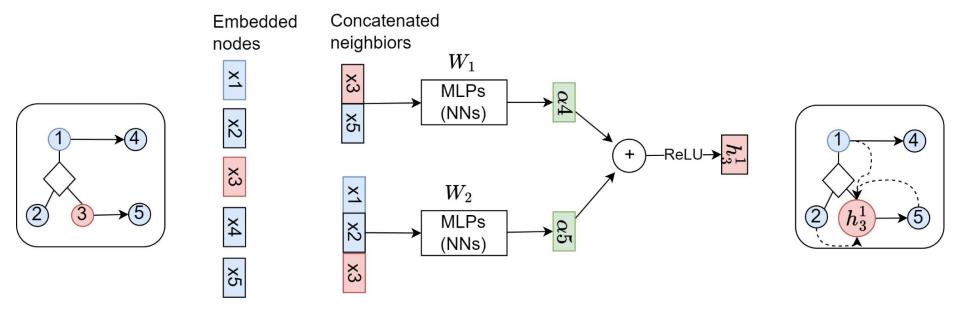


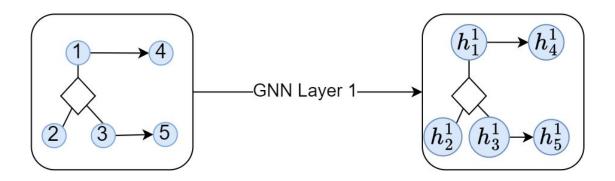


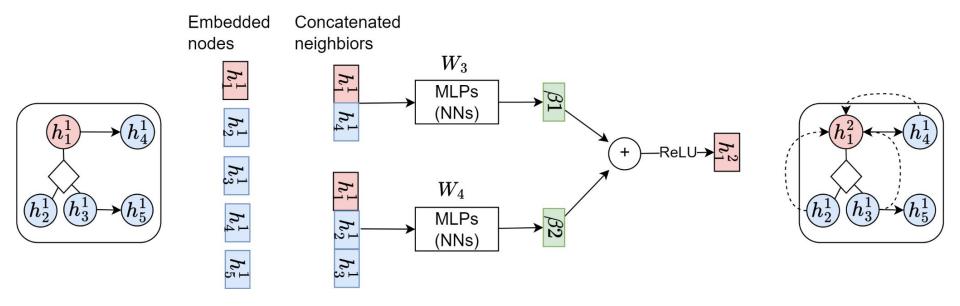


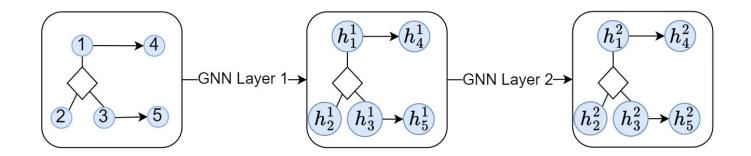












Use predicted MUSes to guide the algorithms

- Prioritize CHCs by using predicted scores of CHCs
 - Use scores alone
 - Combine with original prioritizing scores
 - Add/subtract normalized or ranked scores with coefficient
 - Randomly shift to MUS and original score

Algorithm	Name
	Fixed
8	Random
CEGAR	Score
·-	Rank
9.	R-Plus
	S-Plus
ă.	R-Minus
	S-Minus
	Fixed
٥.	Random
SymEx	Score
Symex	Rank
	R-Plus
25	S-Plus
	R-Minus
	S-Minus
1	Two-queue

Experimental results

Benchmarks from CHC-COMP

Linear LIA proble	ems	Non-linear LIA problems			
8705		8425			
Benchmarks for training	Holdout set Benchmarks for training H		Holdout set		
7834 (90%)	871 (10%)	7579 (90%) 846 (846 (10%)	
UNSAT SAT T/O	Eval. N/A	UNSAT	SAT T/O	Eval. N/A	
1585 4004 2245	383 488	3315	4010 254	488 358	
Train Valid N/A		Train Valid N/A			
782 87 716		1617 180 1518			

Experimental results (Improved percentage)

Benchmark	MUS	Best ranking function (improvement in %)									
Algorithm	data set		of Solved	Problems	Average Time						
	(best count)	Total	SAT	UNSAT	All	Common	SAT	UNSAT			
	Union	R-Plus	R-Plus	R-Minus	R-Plus	S-Plus	S-Minus	Rank			
Linear	(0)	(1.4%)	(2.4%)	(1.0%)	(1.3%)	(19.1%)	(46.5%)	(31.1%)			
CEGAR	Single	Rank R-Plus		Rank	R-Plus	S-Plus	R-Minus	Rank			
enter de la	(3)	(3.6%)	(4.0%)	(8.2%)	(1.9%)	(26.6%)	(57.9%)	(36.3%)			
	Intersection	R-Plus	S-Plus	R-Plus	R-Plus	S-Plus	R-Minus	S-Plus			
955	(4)	(4.1%)	(0.8%)	(9.3%)	(3.1%)	(27.6%)	(45.0%)	(0.0%)			
	Union	Two-Q	S-Plus*	Random	Two-Q	R-Minus	R-Minus	S-Plus			
Linear	(4)	(1.0%)	(0.0%)	(2.0%)	(0.9%)	(12.7%)	(30.2%)	(26.5%)			
SymEx	Single	S-Minus*	S-Plus*	Random	Random	S-Plus	Random	S-Plus			
	(3)	(0.5%)	(0.0%)	(2.0%)	(0.8%)	(12.9%)	(28.4%)	(17.6%)			
	Intersection	S-Plus*	S-Plus*	S-Plus*	S-Plus	Score	Random	R-Plus			
	(5)	(1.0%)	(0.0%)	(2.0%)	(1.3%)	(9.5%)	(28.4%)	(35.8%)			

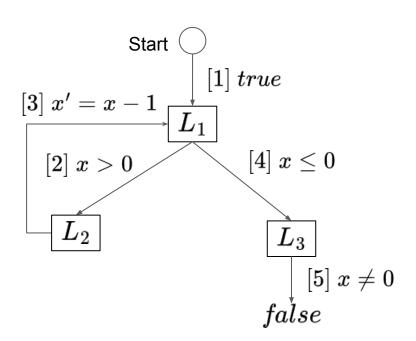
Experimental results (Improved percentage)

Benchmark	MUS	Best ranking function (improvement in %)								
Algorithm	data set	Numbe	er of Solved	Problems	Average Time					
TI gorronn	(best count)	Total	Total SAT UNSAT All		All	Common	SAT	UNSAT		
Non- Linear CEGAR	Union (7)	S-Plus (0.5%)	S-Plus (0.8%)	S-Plus* (0.0%)	S-Plus (7.1%)	R-Minus (20.8%)	Rank (53.5%)	S-Plus (19.4%)		
	Single (1) Intersection	R-Plus (0.2%) R-Plus*	R-Plus (0.4%) S-Plus	R-Plus* (0.0%) S-Plus*	R-Plus (6.6%) R-Plus	S-Plus (18.4%) R-Plus	R-Minus (52.8%) Rank	R-Minus (14.2%) S-Plus		
	(1) Union	(0.0%) Two-Q	(0.5%) S-Minus*	(0.0%) Random	(5.9%)	(20.3%) R-Minus	(45.8%) Score	(16.8%) R-Plus		
Non- Linear	(6)	(6.1%)	(1.6%)	(12.3%)	(13.3%)	(7.3%)	(5.1%)	(19.9%)		
SymEx	Single (3)	Two-Q (6.1%)	Score (1.6%)	Two-Q (12.9%)	Two-Q (12.4%)	Rank (-2.2%)	R-Minus (0.2%)	Two-Q (11.2%)		
	Intersection (3)	Two-Q (6.1%)	S-Plus (1.6%)	Two-Q (12.9%)	Two-Q (12.7%)	S-Minus (0.6%)	Two-Q (1.7%)	S-Plus (5.4%)		

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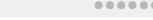
Visualize CHCs with dependency graph

$$egin{array}{lll} [1] \ L_1(x) & \leftarrow true \ [2] \ L_2(x) & \leftarrow L_1(x) \wedge x > 0 \ [3] \ L_1(x') & \leftarrow L_2(x) \wedge x' = x - 1 \ [4] \ L_3(x) & \leftarrow L_1(x) \wedge x \leq 0 \ [5] \ false & \leftarrow L_3(x) \wedge x
eq 0 \end{array}$$



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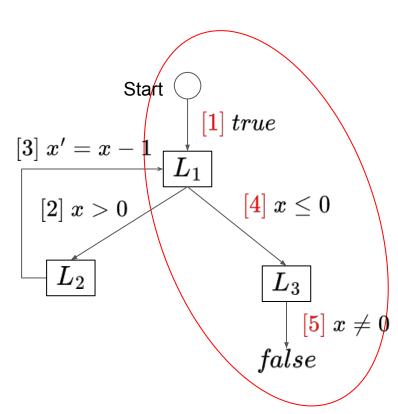
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MUSes of CHCs

$$egin{array}{lll} [1] \ L_1(x) & \leftarrow true \ [2] \ L_2(x) & \leftarrow L_1(x) \wedge x > 0 \ [3] \ L_1(x') & \leftarrow L_2(x) \wedge x' = x - 1 \ [4] \ L_3(x) & \leftarrow L_1(x) \wedge x \leq 0 \ [5] \ false & \leftarrow L_3(x) \wedge x
eq 0 \end{array}$$

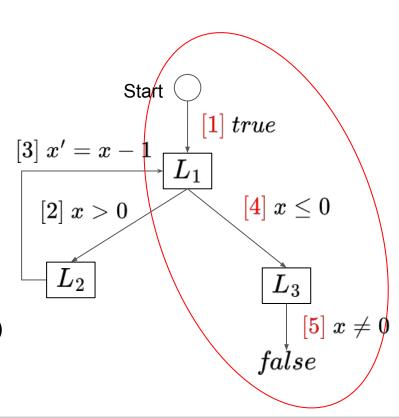
{[1], [4], [5]} is the only MUSes



MUSes of CHCs

$$egin{array}{lll} [1] \ L_1(x) & \leftarrow true \ [2] \ L_2(x) & \leftarrow L_1(x) \wedge x > 0 \ [3] \ L_1(x') & \leftarrow L_2(x) \wedge x' = x - 1 \ [4] \ L_3(x) & \leftarrow L_1(x) \wedge x \leq 0 \ [5] \ false & \leftarrow L_3(x) \wedge x
eq 0 \end{array}$$

- Algorithms
 - Counterexample-guided abstraction refinement (CEGAR)
 - Symbolic execution (Symex)



.

MUSes of CHCs

Score

0.8 [1]
$$L_1(x) \leftarrow true$$

0.2 [2]
$$L_2(x) \leftarrow L_1(x) \land x > 0$$

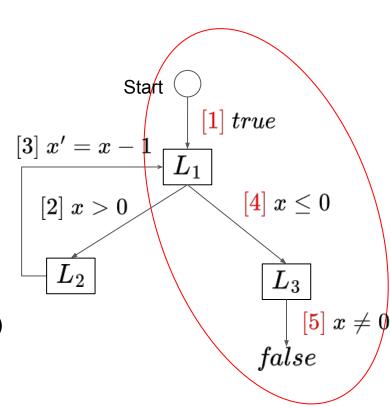
0.1
$$[3]$$
 $L_1(x')$ $\leftarrow L_2(x) \wedge x' = x-1$

0.75 [4]
$$L_3(x) \leftarrow L_1(x) \land x \leq 0$$

0.6 [5]
$$false \leftarrow L_3(x) \land x \neq 0$$

Algorithms

- Counterexample-guided abstraction refinement (CEGAR)
- Symbolic execution (Symex)



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Framework overview

Extract program features by graph neural networks from CHC's graph representation



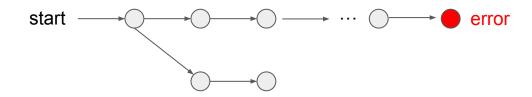
Motivation & Challenges

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Five proxy tasks from simple to difficult

- 1. Predict if an graph node is an argument of relation symbol
- 2. Predict the number of occurrence of the relation symbols in all clauses

- 3. Predict if a relation symbol is in a cycle
- 4. Predict the existence of argument bound
- 5. Predict the clause membership in all/some minimal unsat cores



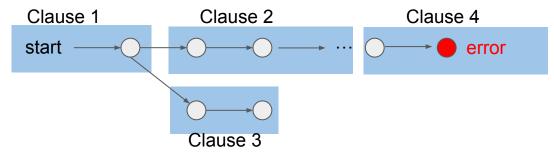
Motivation & Challenges

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Five proxy tasks from simple to difficult

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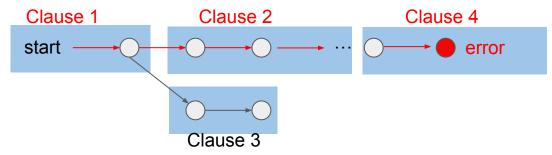
Motivation & Challenges

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Five proxy tasks from simple to difficult

- 1. Predict if an graph node is an argument of relation symbol
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Motivation & Challenges

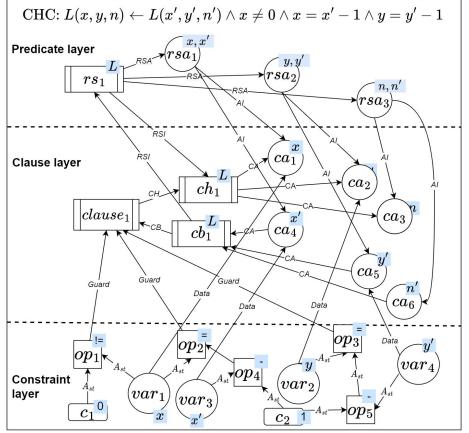
Framework overview

Two graph representations for CHCs



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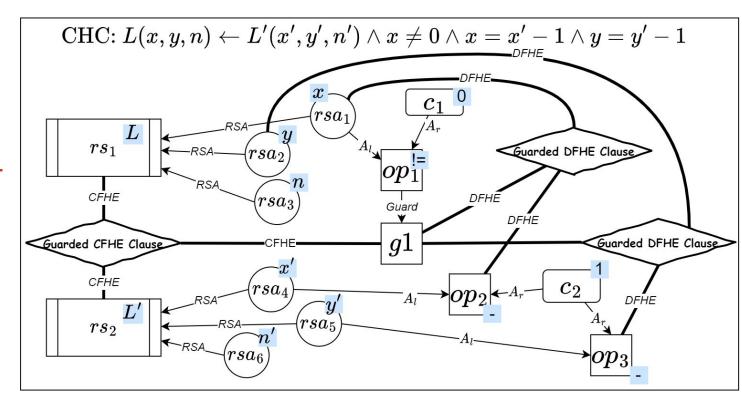
Constraint graph (CG)



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Control- and dataflow hypergraph (CDHG)



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Framework overview

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Motivation & Challenges

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Relational Hypergraph Neural Network (R-HyGNN)

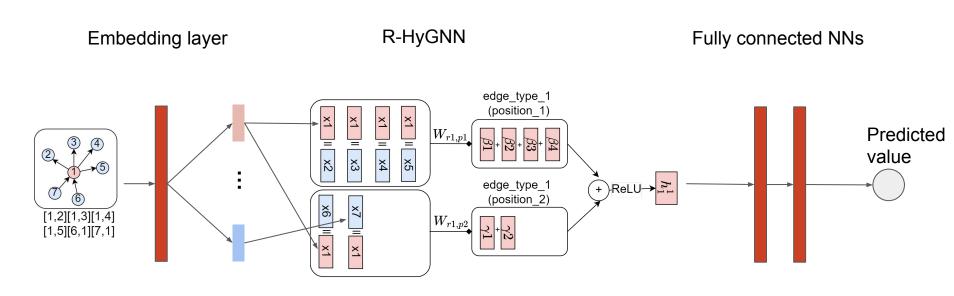


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Training model

Motivation & Challenges

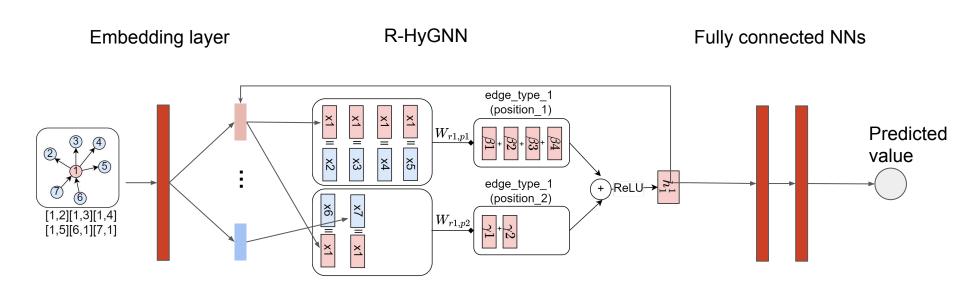


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Training model

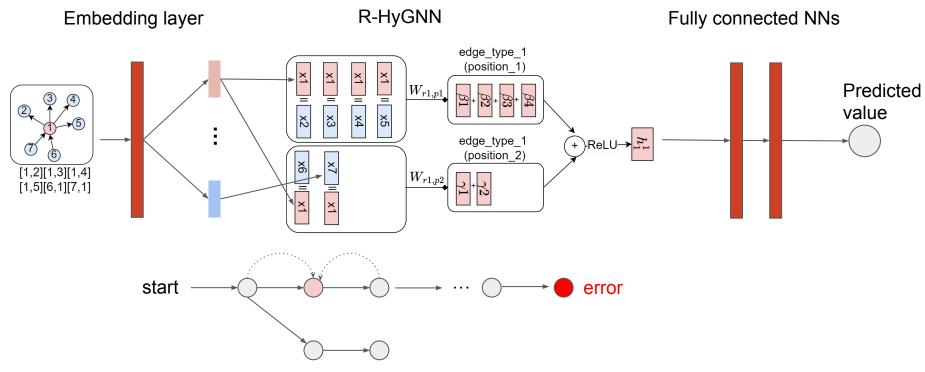
Motivation & Challenges



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Motivation & Challenges

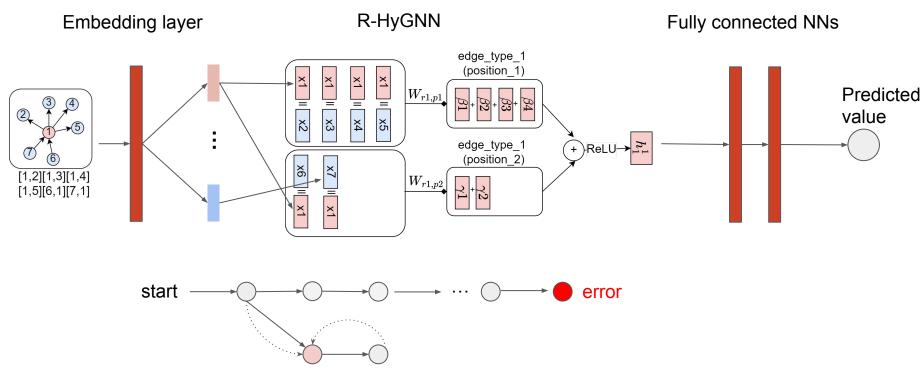
Training model



Training model

Motivation & Challenges

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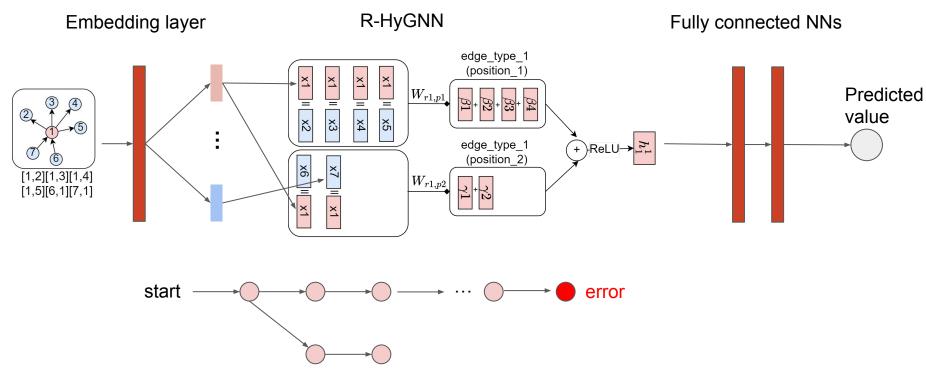


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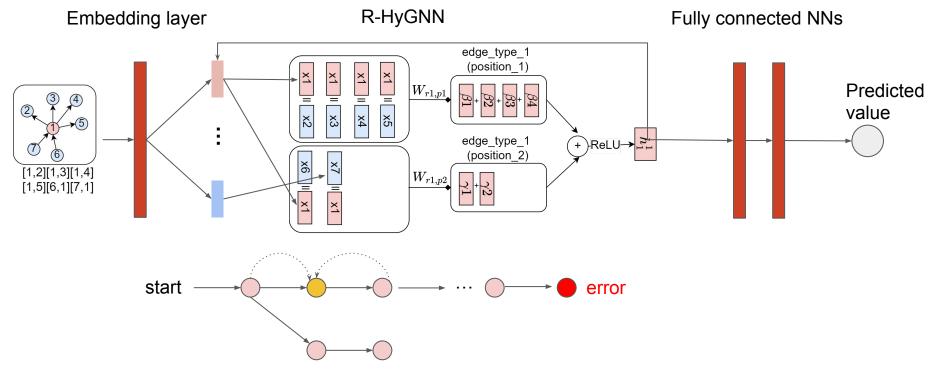
Motivation & Challenges

Training model





Motivation & Challenges



Framework overview

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Motivation & Challenges



Evaluation on five proxy tasks

Dominate distribution

Motivation & Challenges

				CC	;		CDHG				
Task	Files	T P	+	-	Acc.	Dom.	+	-	Acc.	Dom.	
1	1121	+	93863	0	100%	95.1%	142598	0	99.9%	72.8%	
I.	1121	-	0	1835971	100%		10	381445			
3	1115	+	3204	133	96.1%	70.1%	8262	58	99.6%	50.7%	
3	3 1113	-	301	7493			15	8523			
4 (a)	1 (0)	+	13685	5264	91.2%	79.7%	30845	4557	94.3%	75.2%	
4 (a)	1028	-	2928	71986			3566	103630			
4 (b)	1028	+	18888	4792	91.4%	74.8%	41539	4360	94.3%	67.8%	
4 (0)		-	3291	66892			3715	92984			
5 (a)	386	+	1048	281	95.0%	84.7%	1230	206	96.9%	86.4%	
3 (a) 300	360	E	154	7163		04.770	121	9036			
5 (b)	383	+	3030	558	84.6%	53.1%	3383	481	90.6%	54.8%	
3 (0)	303	-	622	3428			323	4361			

Evaluation on five proxy tasks

Motivation & Challenges

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Accuracy is higher than dominate distribution

				CC	j		CDHG				
Task	Files	T P	+	-	Acc.	Dom.	+	-	Acc.	Dom.	
1	1121	+	93863	0	100%	95.1%	142598	0	99.9%	72.8%	
<u>.</u>	1121	-	0	1835971	10076		10	381445			
3	1115	+ 3204 133	96.1%	96.1% 70.1%	8262	58	99.6%	50.7%			
3	3 1113	-	301	7493	90.176	70.176	15	8523	77.070	30.770	
4 (a)	1(2)	+	13685	5264	91.2%	79.7%	30845	4557	94.3%	75.2%	
4 (a)	1028	-	2928	71986			3566	103630			
4 (b)	1028	+	18888	4792	91.4%	74.8%	41539	4360	94.3%	67.8%	
4 (0)		-	3291	66892			3715	92984			
5 (a)	386	+	1048	281	95.0%	84.7%	1230	206	96.9%	86.4%	
3 (a) 300	360	i e	154	7163		04.770	121	9036			
5 (b)	383	+	3030	558	84.6%	53.1%	3383	481	90.6%	54.8%	
5 (B)	303	-	622	3428			323	4361			

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Motivation & Challenges

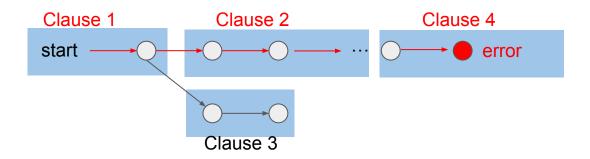
Evaluation on five proxy tasks (Task 5 results)

- Task 5: Predict the clause membership in (a) some and (b) all of the minimal unsat cores
- Task 5 (a) 1155, 386, 386 problems for train, valid, and test respectively

			CG				CDHG				
Task	Files	T P	+	-	Acc.	Dom.	+	-	Acc.	Dom.	
1	1121	+	93863	0	100%	95.1%	142598	0	99.9%	72.8%	
ı	1121	-	0	1835971	10076		10	381445			
3	1115	+	3204	133	96.1%	70.1%	8262	58	99.6%	50.7%	
3	1113	-	301	7493			15	8523			
4 (a)		+	13685	5264	91.2%	79.7%	30845	4557	94.3%	75.2%	
4 (a)	1028	-	2928	71986			3566	103630			
4 (b)	1028	+	18888	4792	91.4%	74.8%	41539	4360	94.3%	67.8%	
4 (0)		-	3291	66892			3715	92984			
5 (a)	386	+	1048	281	95.0%	84.7%	1230	206	96.9%	86.4%	
3 (a) 3	300	-	154	7163		04.770	121	9036			
5 (b)	383	+	3030	558	84.6%	53.1%	3383	481	90.6%	54.8%	
5 (b)	303	-	622	3428			323	4361			

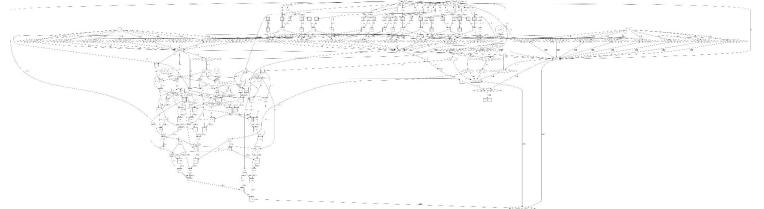
Predict the clause membership in minimal unsat cores

- Simple patterns
 - Clauses close to the assertions are likely in the minimal unsat cores
- Intricate patterns
 - Perfectly predict the clause membership of minimal unsat cores in the case that contain 290 clauses.

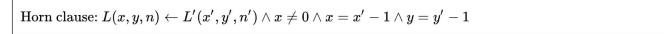


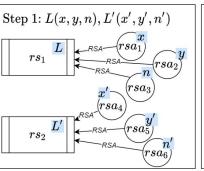
Predict the minimal unsatisfiable cores

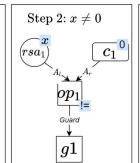
- Simple patterns
 - Clauses close to the assertions are likely in the minimal unsatisfiable cores
- Intricate patterns
 - In the verification problem that contain 290 clauses, the model can perfectly predict the clause membership in minimal unsat cores.

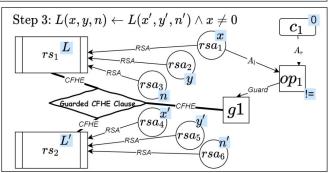


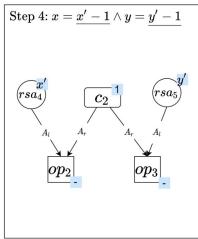
Control- and dataflow hypergraph (CDHG)

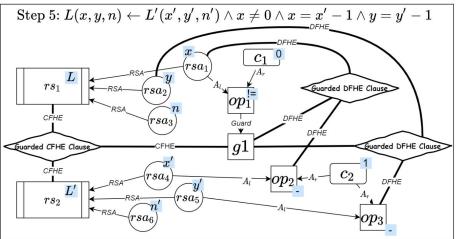




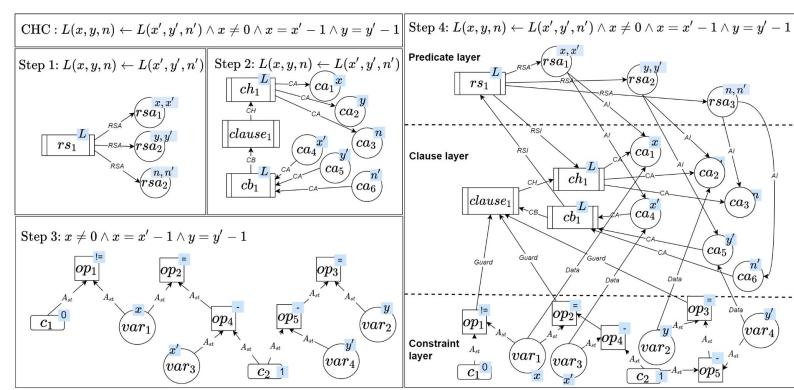


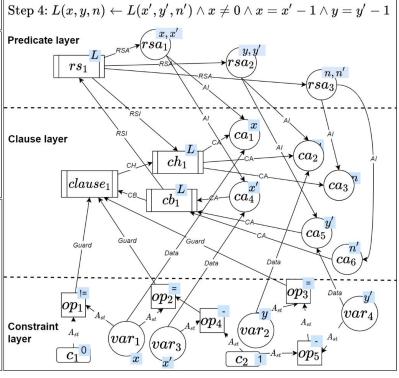






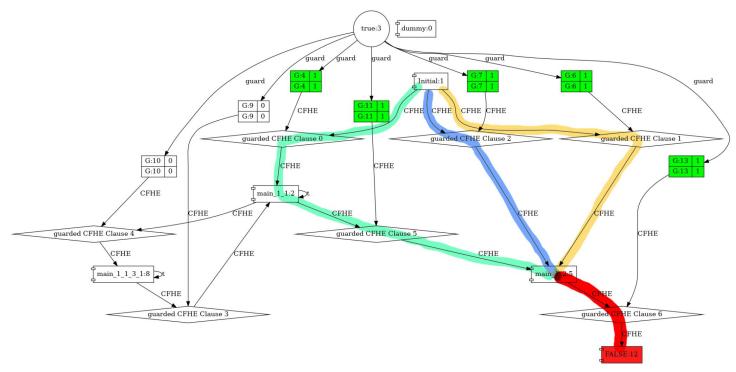
Constraint graph (CG)



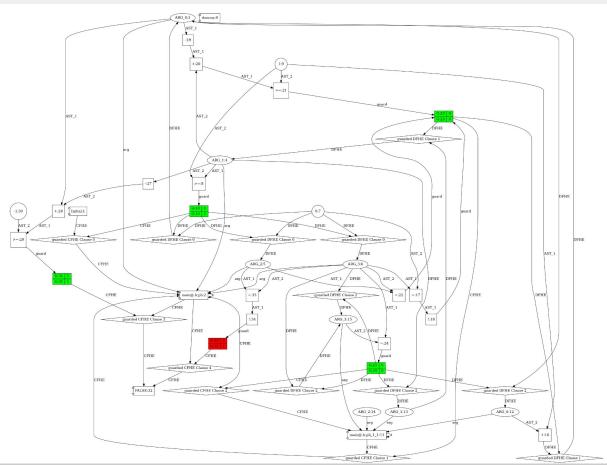


Predict the minimum unsatisfiable cores

Three minimum unsatisfiable cores



Technique details

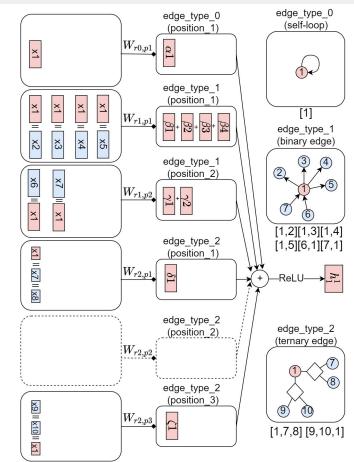


R-HyGNN

The updating rule for node representation in time step t:

$$h_v^t = \text{ReLU}(\sum_{r \in R} \sum_{p \in P_r} \sum_{e \in E_v^{r,p}} W_{r,p}^t \cdot ||[h_u^{t-1} \mid u \in e]),$$

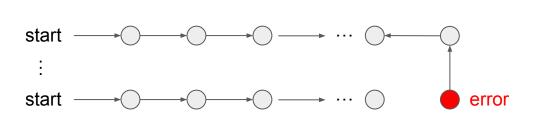
where $\|\{\cdot\}$ denotes concatenation of all elements in a set, $r \in R = \{r_i \mid i \in \mathbb{N}\}$ is the set of edge types (relations), $p \in P_r = \{p_j \mid j \in \mathbb{N}\}$ is the set of node positions under edge type $r, W_{r,p}^t$ denotes learnable parameters when the node is in the pth position with edge type r, and $e \in E_v^{r,p}$ is the set of hyperedges of type r in the graph in which node v appears in position p, where e is a list of nodes.

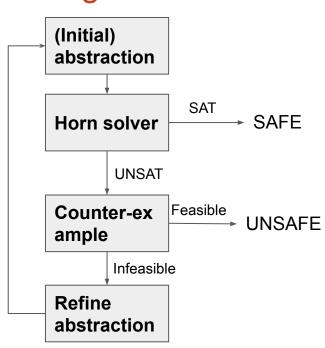


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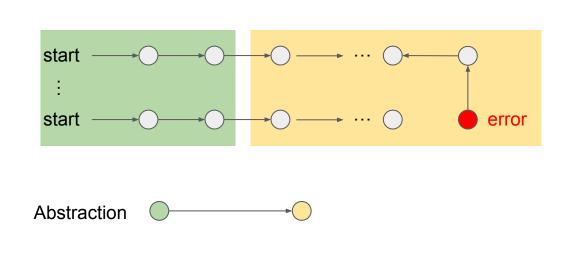
Motivation & Background

Solving CHCs by counter-example guided abstraction refinement (CEGAR) based model checking



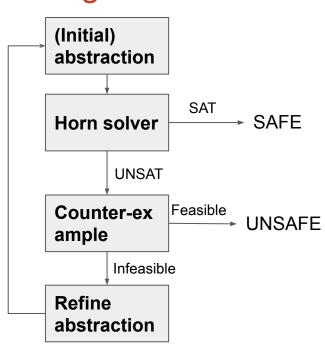


Solving CHCs by counter-example guided abstraction refinement (CEGAR) based model checking

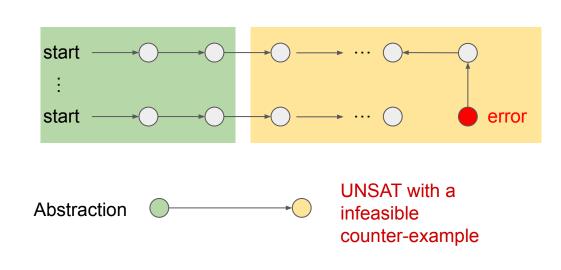


Motivation & Background

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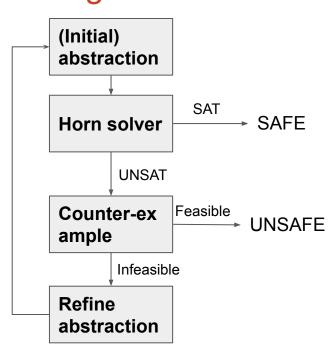


Solving CHCs by counter-example guided abstraction refinement (CEGAR) based model checking



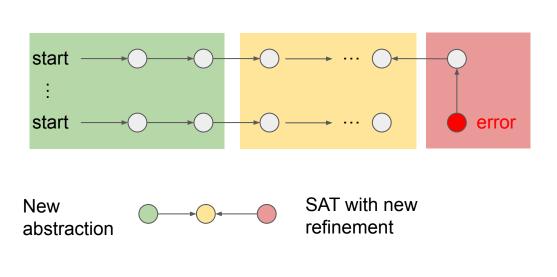
Motivation & Background

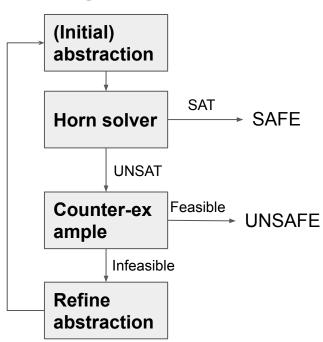
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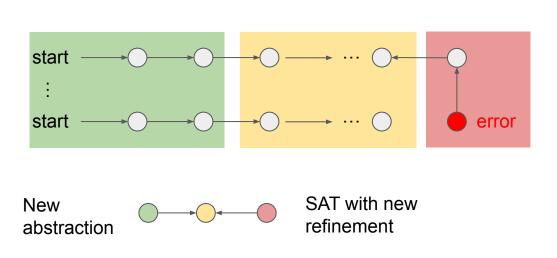
Motivation & Background

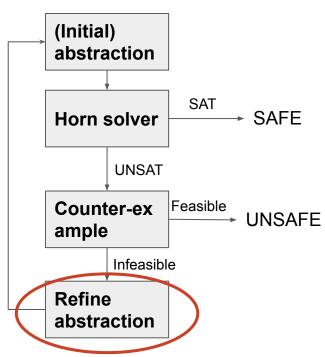
Solving CHCs by counter-example guided abstraction refinement (CEGAR) based model checking





Solving CHCs by counter-example guided abstraction refinement (CEGAR) based model checking





Motivating Examples (Theorem Proving)

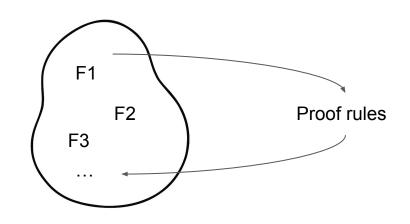
$$ext{Prove}: orall x. P(x)
ightarrow Q(x) \wedge orall x. Q(x)
ightarrow R(x) \wedge P(a) \Longrightarrow R(a)$$

F1:
$$\forall x. P(x) \rightarrow Q(x)$$

F2:
$$\forall x. Q(x) \rightarrow R(x)$$

F3: P(a)

Goal: R(a)



Motivating Examples (Theorem Proving)

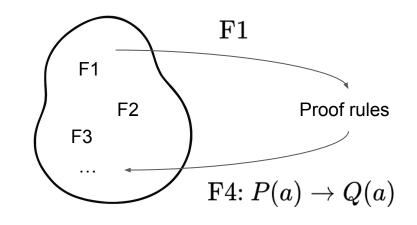
$$\text{Prove}: \forall x.\, P(x) \rightarrow Q(x) \land \forall x.\, Q(x) \rightarrow R(x) \land P(a) \Longrightarrow R(a)$$

F1:
$$\forall x. P(x) \rightarrow Q(x)$$

F2:
$$\forall x. Q(x) \rightarrow R(x)$$

F3: P(a)

Goal: R(a)



Motivating Examples (Theorem Proving)

$$ext{Prove}: orall x. P(x)
ightarrow Q(x) \wedge orall x. Q(x)
ightarrow R(x) \wedge P(a) \Longrightarrow R(a)$$

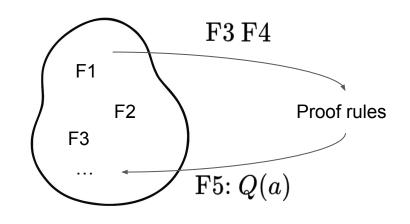
F1:
$$\forall x. P(x) \rightarrow Q(x)$$

F2:
$$\forall x. Q(x) \rightarrow R(x)$$

F3: P(a)

F4: $P(a) \rightarrow Q(a)$

Goal: R(a)



Rank clauses before solving

$$egin{array}{lll} C1: \ L_1(x) & \leftarrow true \ C2: \ L_2(x) & \leftarrow L_1(x) \wedge x > 0 \ C3: \ L_1(x') & \leftarrow L_2(x) \wedge x' = x - 1 \ C4: \ L_3(x) & \leftarrow L_1(x) \wedge x \leq 0 \ C5: \ false & \leftarrow L_3(x) \wedge x = 0 \end{array}$$

Rank clauses before solving

$$egin{array}{lll} C1: L_1(x) & \leftarrow true \ C2: L_2(x) & \leftarrow L_1(x) \wedge x > 0 \ C3: L_1(x') & \leftarrow L_2(x) \wedge x' = x - 1 \ C4: L_3(x) & \leftarrow L_1(x) \wedge x \leq 0 & \longrightarrow C6: L_3 \leftarrow true \wedge x \leq 0 \ C5: false & \leftarrow L_3(x) \wedge x = 0 \ \end{array}$$

Rank clauses before solving

$$egin{array}{lll} C1: L_1(x) & \leftarrow true \ C2: L_2(x) & \leftarrow L_1(x) \wedge x > 0 \ C3: L_1(x') & \leftarrow L_2(x) \wedge x' = x - 1 \ C4: L_3(x) & \leftarrow L_1(x) \wedge x \leq 0 & \longrightarrow C6: L_3 \leftarrow true \wedge x \leq 0 \ C5: false & \leftarrow L_3(x) \wedge x = 0 & \downarrow \ \end{array}$$

Rank clauses before solving

$$egin{array}{lll} C1: \ L_1(x) & \leftarrow true \ C2: \ L_2(x) & \leftarrow L_1(x) \wedge x > 0 \ C3: \ L_1(x') & \leftarrow L_2(x) \wedge x' = x - 1 \ C4: \ L_3(x) & \leftarrow L_1(x) \wedge x \leq 0 \end{array}$$

$$C5: false \leftarrow L_3(x) \land x = 0$$

Proof rules C3 $-C6:L_3 \leftarrow \hat{t}rue \wedge x \leq 0$

$$C7: false \leftarrow x \leq 0 \land x = 0$$